scal.cs.uni-salzburg.at

multicore-scalable concurrent data structures

scalloc.cs.uni-salzburg.at

multicore-scalable concurrent allocator

selfie.cs.uni-salzburg.at

self-referential systems software for teaching

Scal, Scalloc, and Selfie

Christoph Kirsch, University of Salzburg, Austria

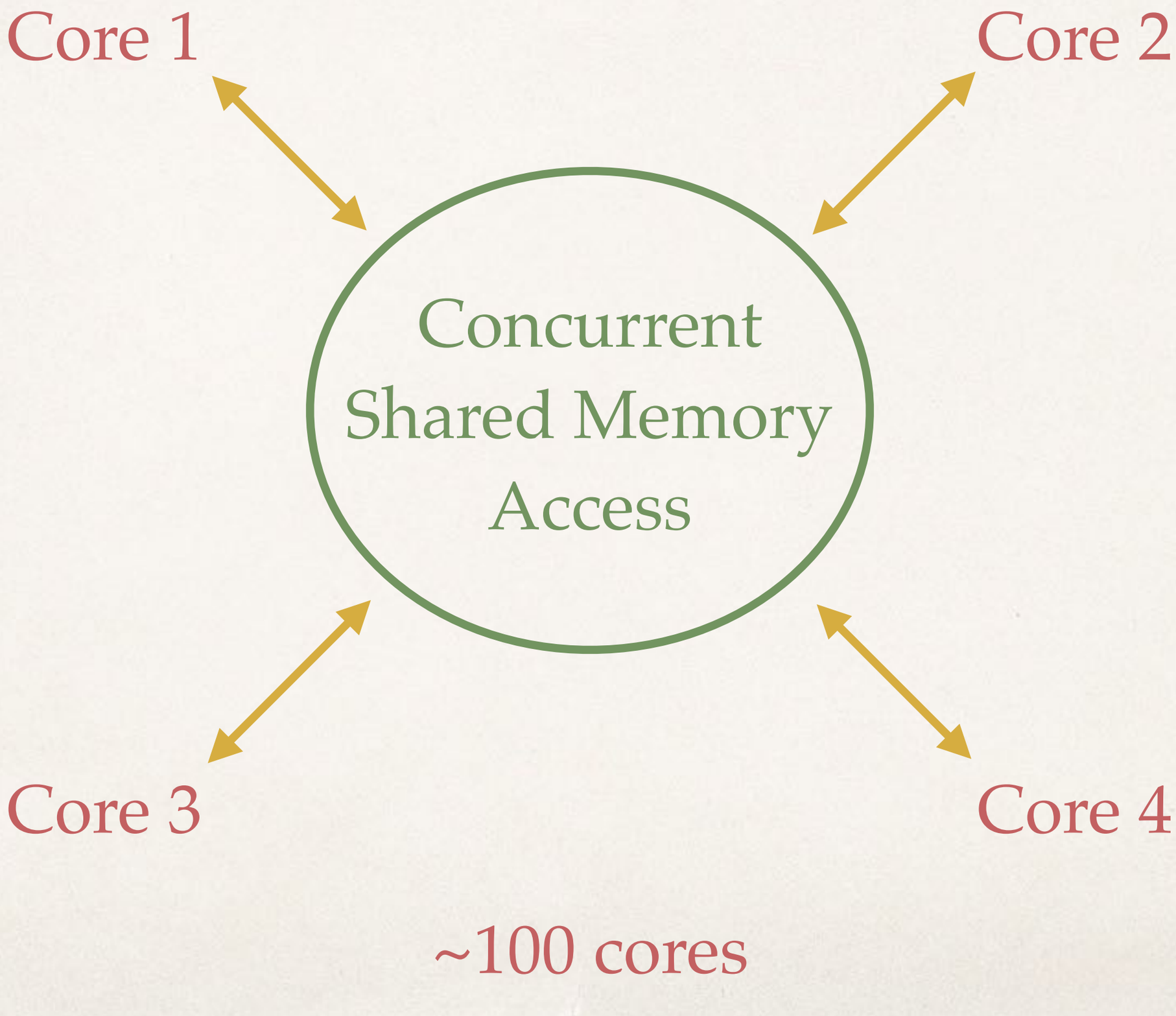
University of Cambridge, September 2017

Joint Work

- Martin Aigner
- Christian Barthel
- Mike Dodds
- Andreas Haas
- Thomas Henzinger
- Andreas Holzer
- Thomas Hütter
- Michael Lippautz
- Alexander Miller
- Simone Oblasser
- Hannes Payer
- Mario Preishuber
- Ana Sokolova
- Ali Szegin

How do we exchange data among increasingly many cores on a shared memory machine such that performance still increases with the number of cores?

–The Multicore Scalability Challenge



~100 cores

How do we allocate and deallocate shared memory with increasingly many cores such that performance increases with the number of cores while memory consumption stays low?

–Multicore Shared Memory Allocation Problem

~1TB memory

Core 1

Core 2

Concurrent
Shared Memory
Allocation

Core 3

Core 4

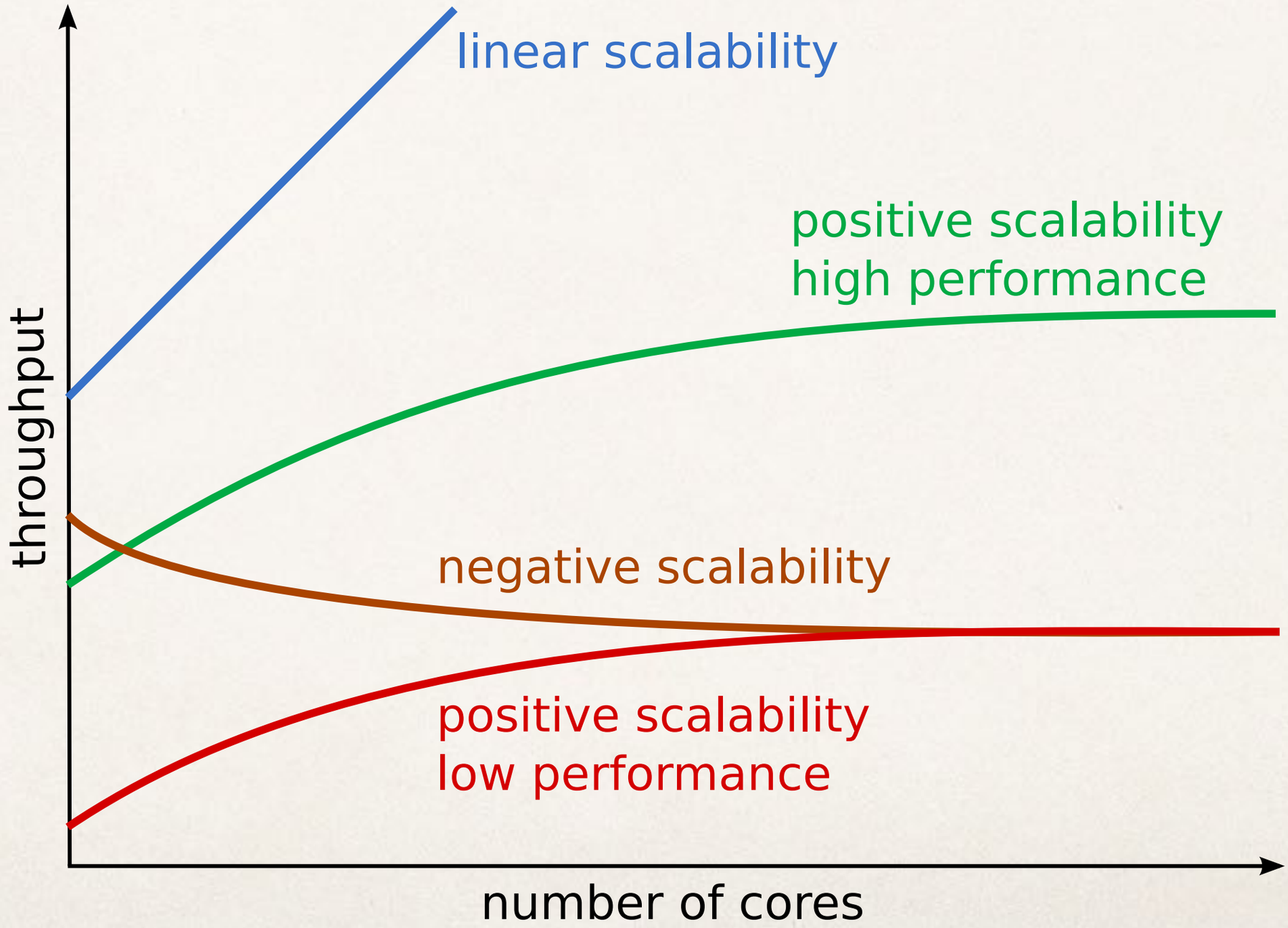
~100 cores

How do we teach computer science to students <u>not necessarily majoring in computer science</u> but who anyway <u>code every day</u>?

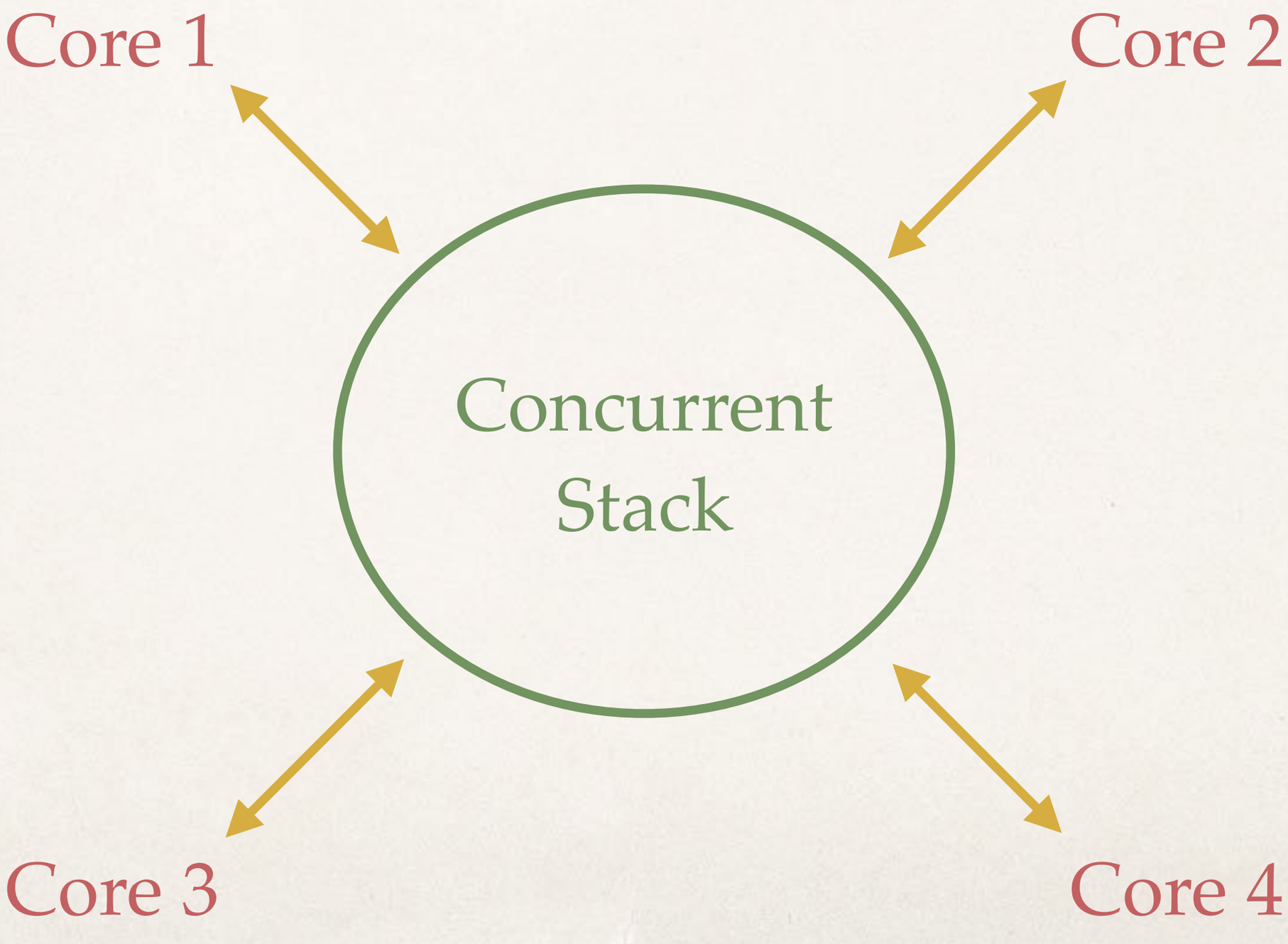
–The Computer Science Education Challenge


OPERA BAR

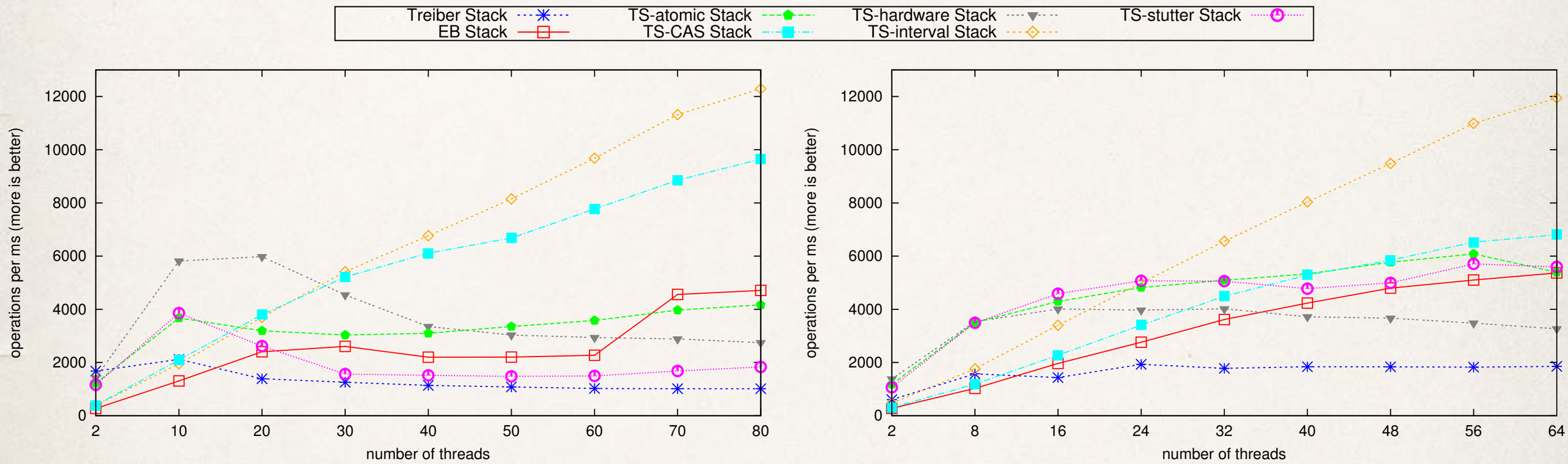
The Multicore Scalability Challenge



Timestamped (TS) Stack [POPL15]



Timestamped (TS) Stack [POPL15]



(a) Producer-consumer benchmark, 40-core machine.

(b) Producer-consumer benchmark, 64-core machine.

Elimination Through Shared Timestamps

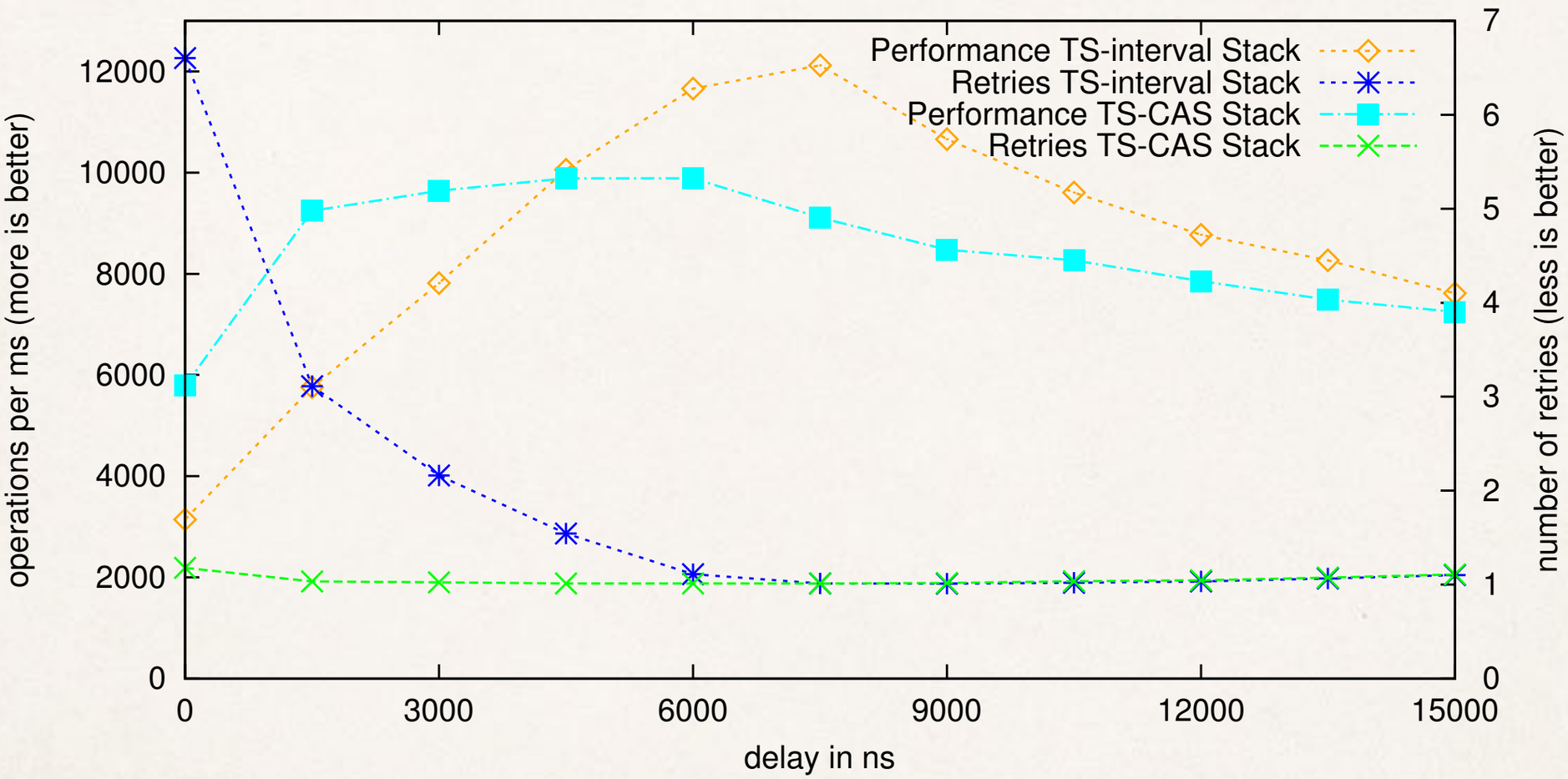
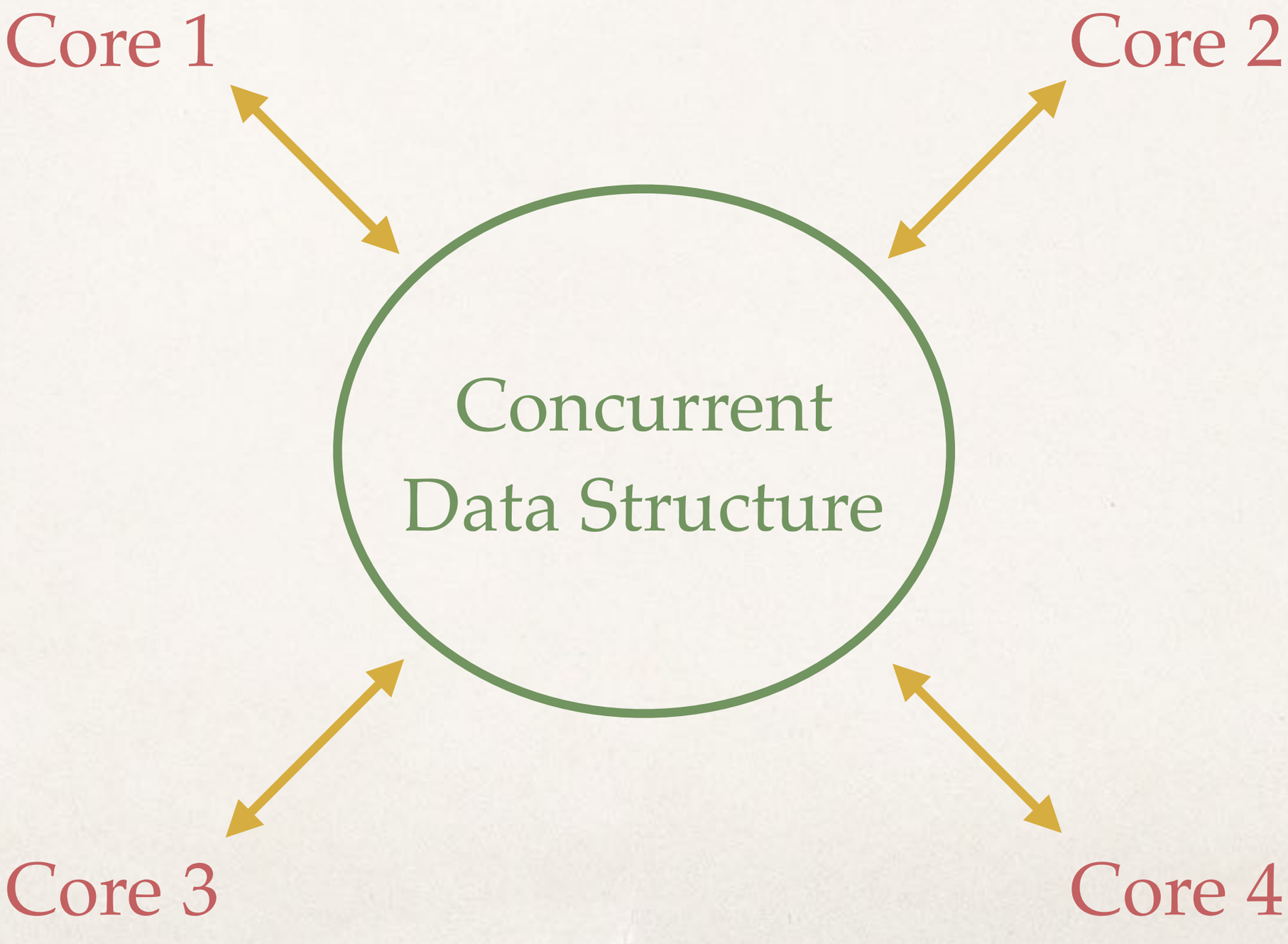


Figure 6: High-contention producer-consumer benchmark using TS-interval and TS-CAS timestamping with increasing delay on the 40-core machine, exercising 40 producers and 40 consumers.

The order in which concurrently pushed elements appear on the TS stack is only determined when they are popped off the TS stack, even if they had been on the TS stack for, say, a year.

–Showing that the TS Stack is a stack is hard hence POPL

Concurrent Data Structures: scal.cs.uni-salzburg.at [POPL13, CF13, POPL15, NETYS15]



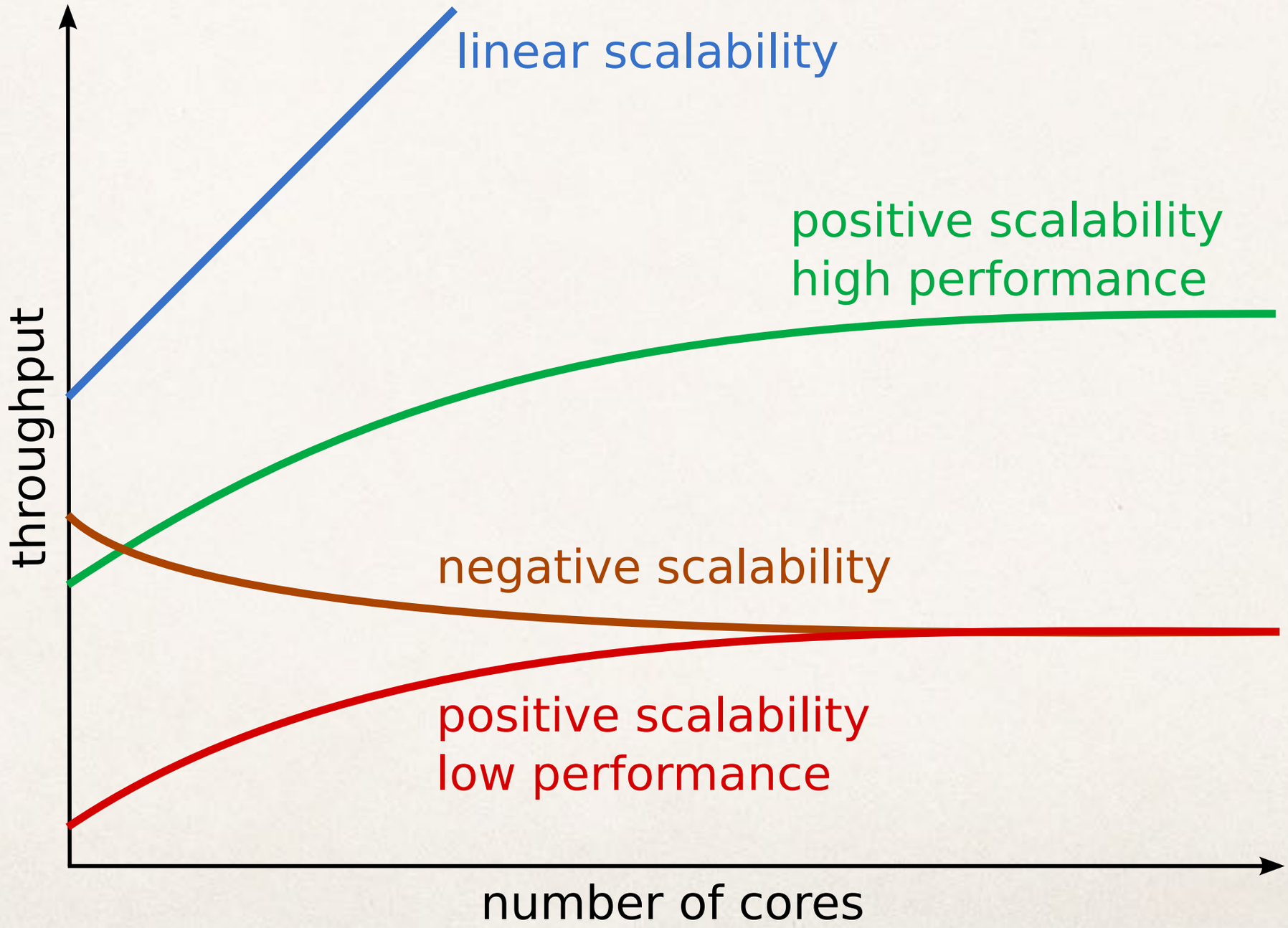
Scal: A Benchmarking Suite for Concurrent Data Structures [NETYS15]

Name	Semantics	Year	Ref
Lock-based Singly-linked	strict queue	1968	[1]
Michael Scott (MS) Queue	strict queue	1996	[2]
Flat Combining Queue	strict queue	2010	[3]
Wait-free Queue	strict queue	2012	[4]
Linked Cyclic Ring Queue	strict queue	2013	[5]
Timestamped (TS) Queue	strict queue	2015	[6]
Cooperative TS Queue	strict queue	2015	[7]
Segment Queue	k-relaxed queue	2010	[8]
Random Dequeue (RD)	k-relaxed queue	2010	[8]
Bounded Size k-FIFO	k-relaxed queue, pool	2013	[9]
Unbounded Size k-FIFO	k-relaxed queue, pool	2013	[9]
b-RR Distributed Queue	k-relaxed queue, pool	2013	[10]
Least-Recently-Used (LRU)	k-relaxed queue, pool	2013	[10]
Locally Linearizable DQ	locally linearizable	2015	[11]
Locally Linearizable k-FIFO	locally linearizable	2015	[11]
Relaxed TS Queue	quiescently consistent	2015	[7]
Lock-based Singly-linked	strict stack	1968	[1]
Treiber Stack	strict stack	1986	[12]
Elimination-backoff Stack	strict stack	2004	[13]
Timestamped (TS) Stack	strict stack	2015	[6]
k-Stack	k-relaxed stack	2013	[14]
b-RR Distributed Stack (DS)	k-relaxed stack, pool	2013	[10]
Least-Recently-Used (LRU)	k-relaxed stack, pool	2013	[10]
Locally Linearizable DS	locally linearizable	2015	[11]
Locally Linearizable k-Stack	locally linearizable	2015	[11]
Timestamped (TS) Deque	strict deque	2015	[7]
d-RA DQ and DS	strict pool	2013	[10]

How do we allocate and deallocate shared memory with increasingly many cores such that performance increases with the number of cores while memory consumption stays low?

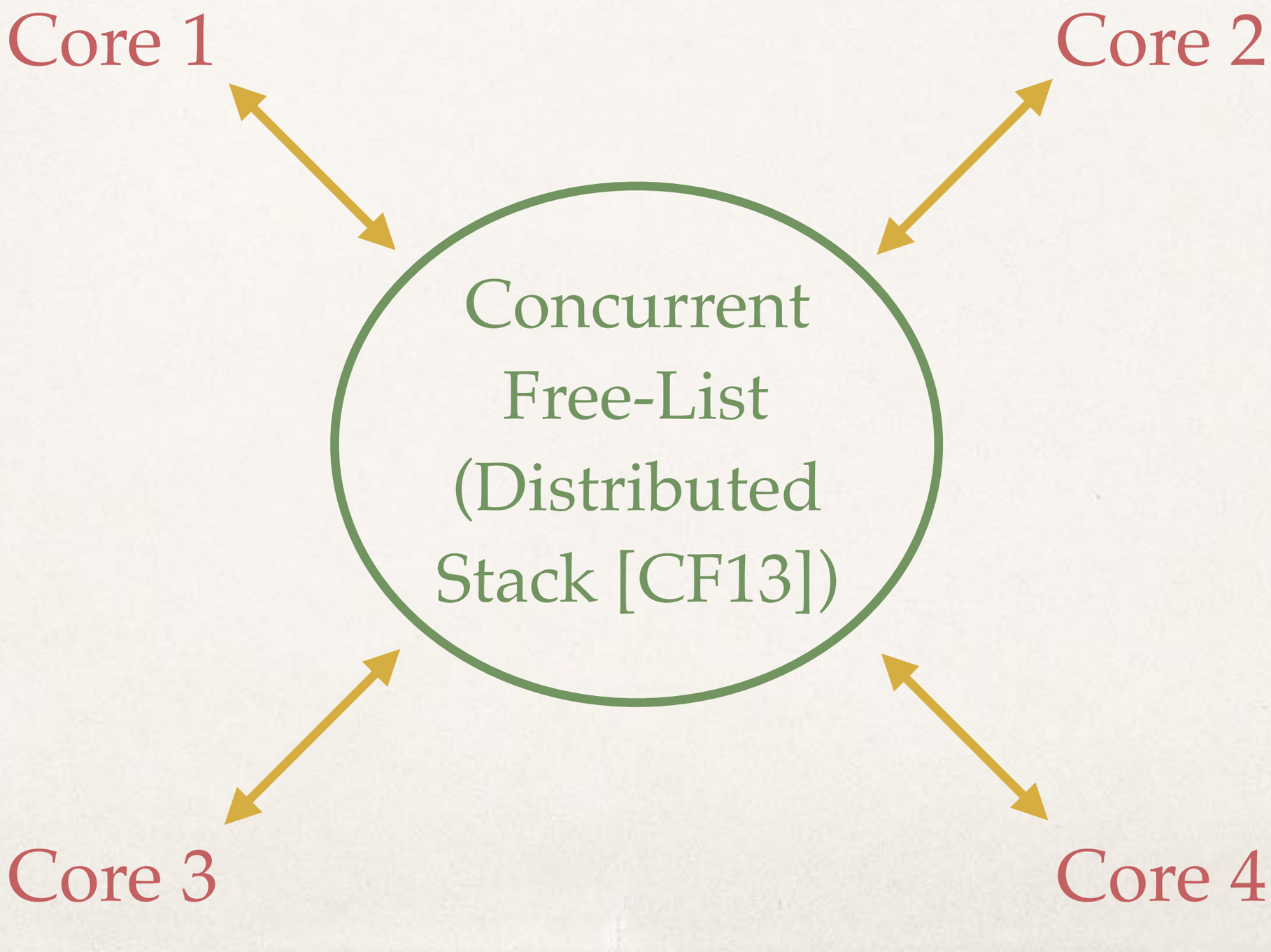
–Multicore Shared Memory Allocation Problem

Multicore Memory Allocation Problem

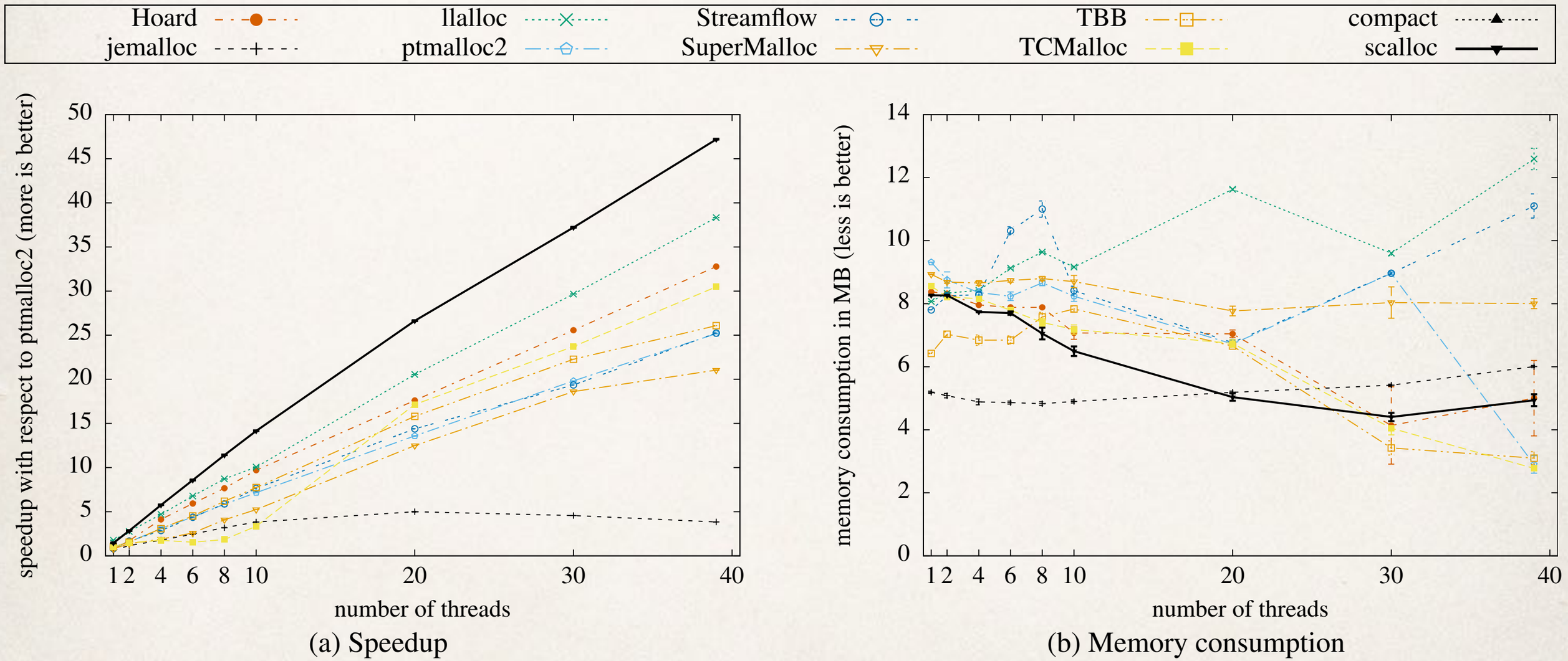


Scalloc: Concurrent Memory Allocator

scalloc.cs.uni-salzburg.at [OOPSLA15]



Local Allocation & Deallocation

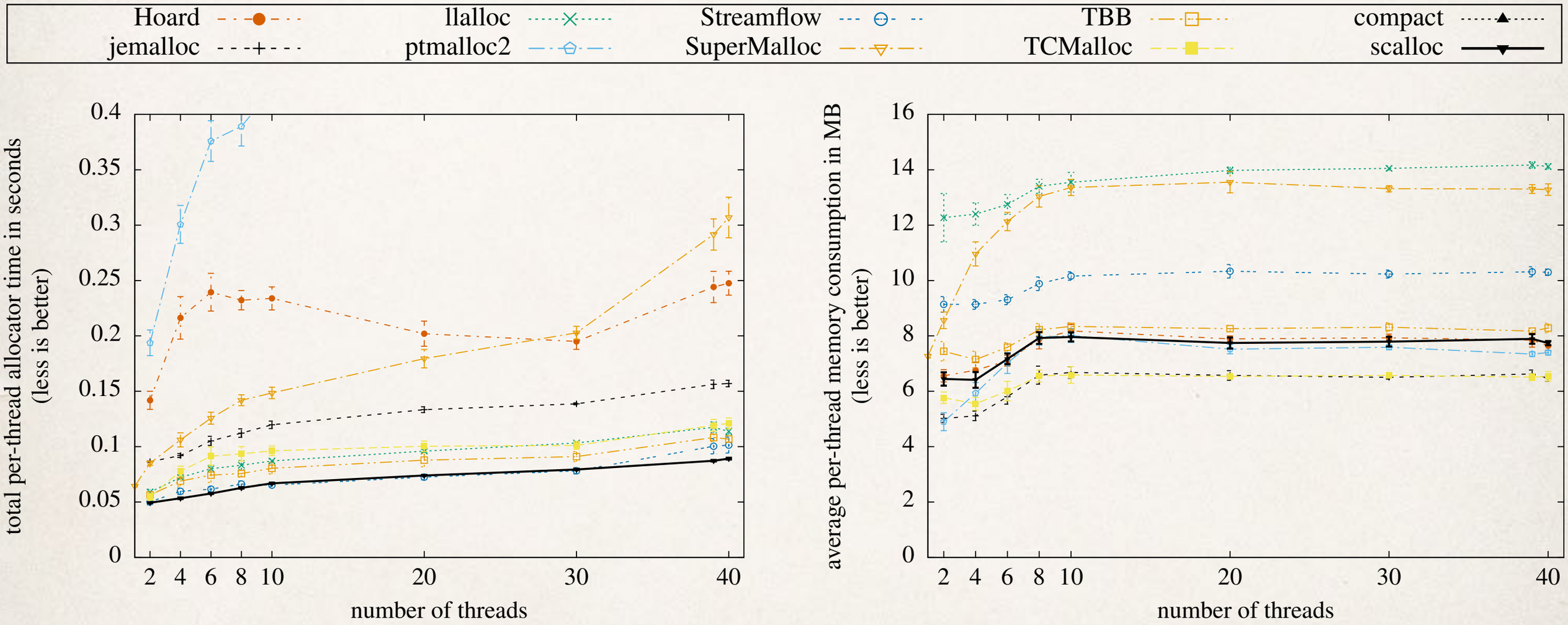


(a) Speedup

(b) Memory consumption

Figure 6: Thread-local workload: Threadtest benchmark

“Remote” Deallocation



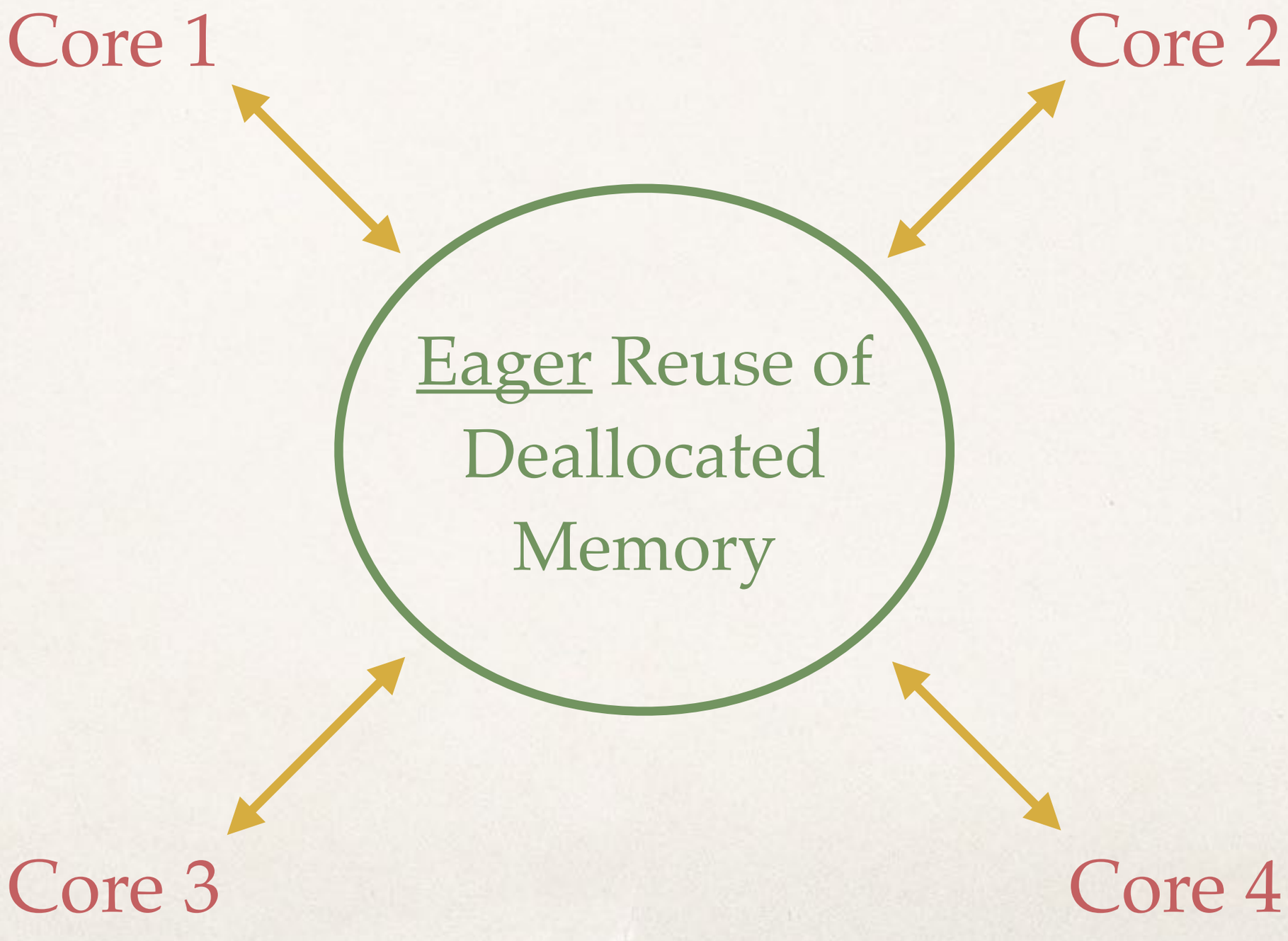
(a) Total per-thread allocator time

(b) Average per-thread memory consumption

Figure 9: Temporal and spatial performance for the producer-consumer experiment

Scalloc: Concurrent Memory Allocator

scalloc.cs.uni-salzburg.at [OOPSLA15]



Virtual Spans: 64-bit Address Space

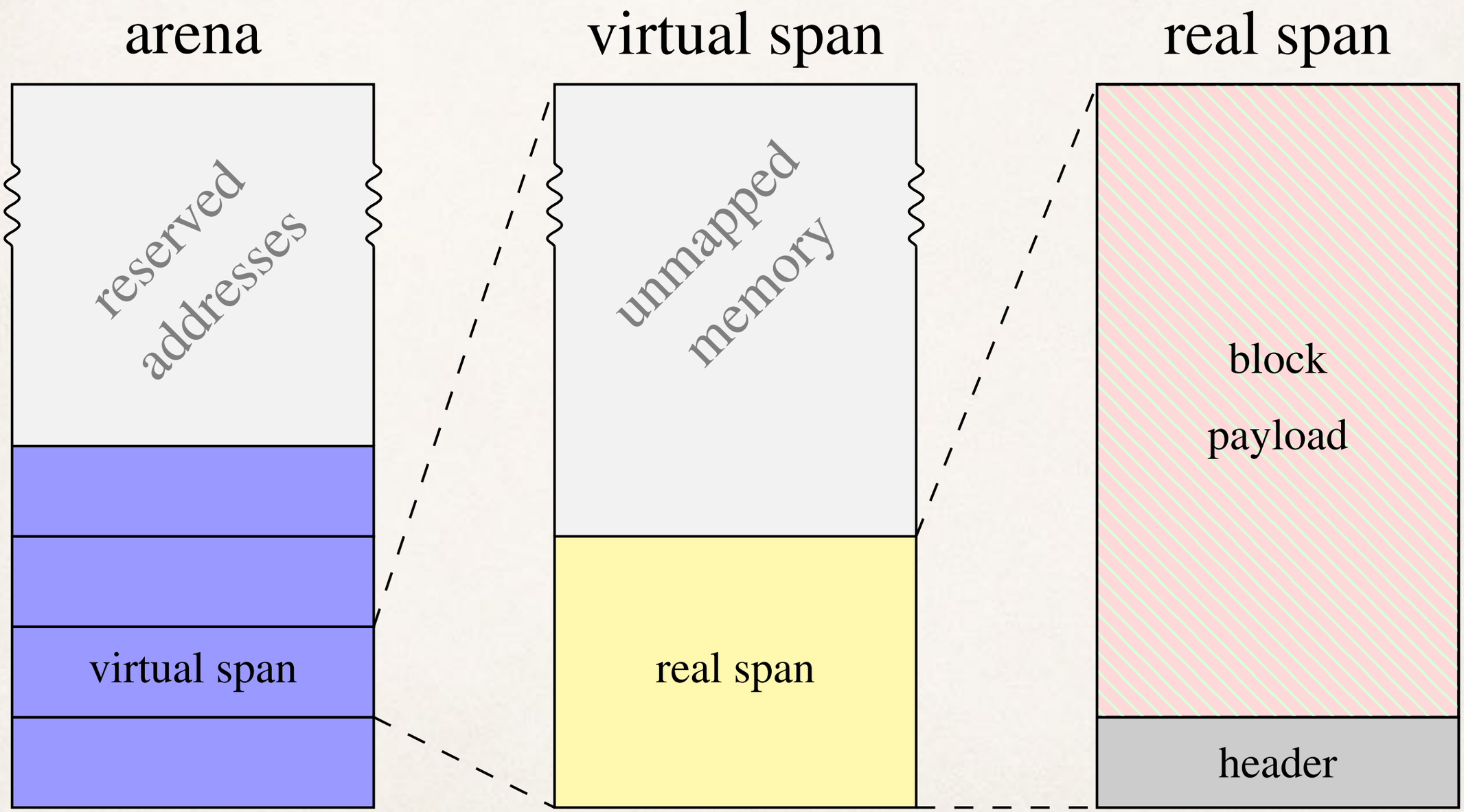
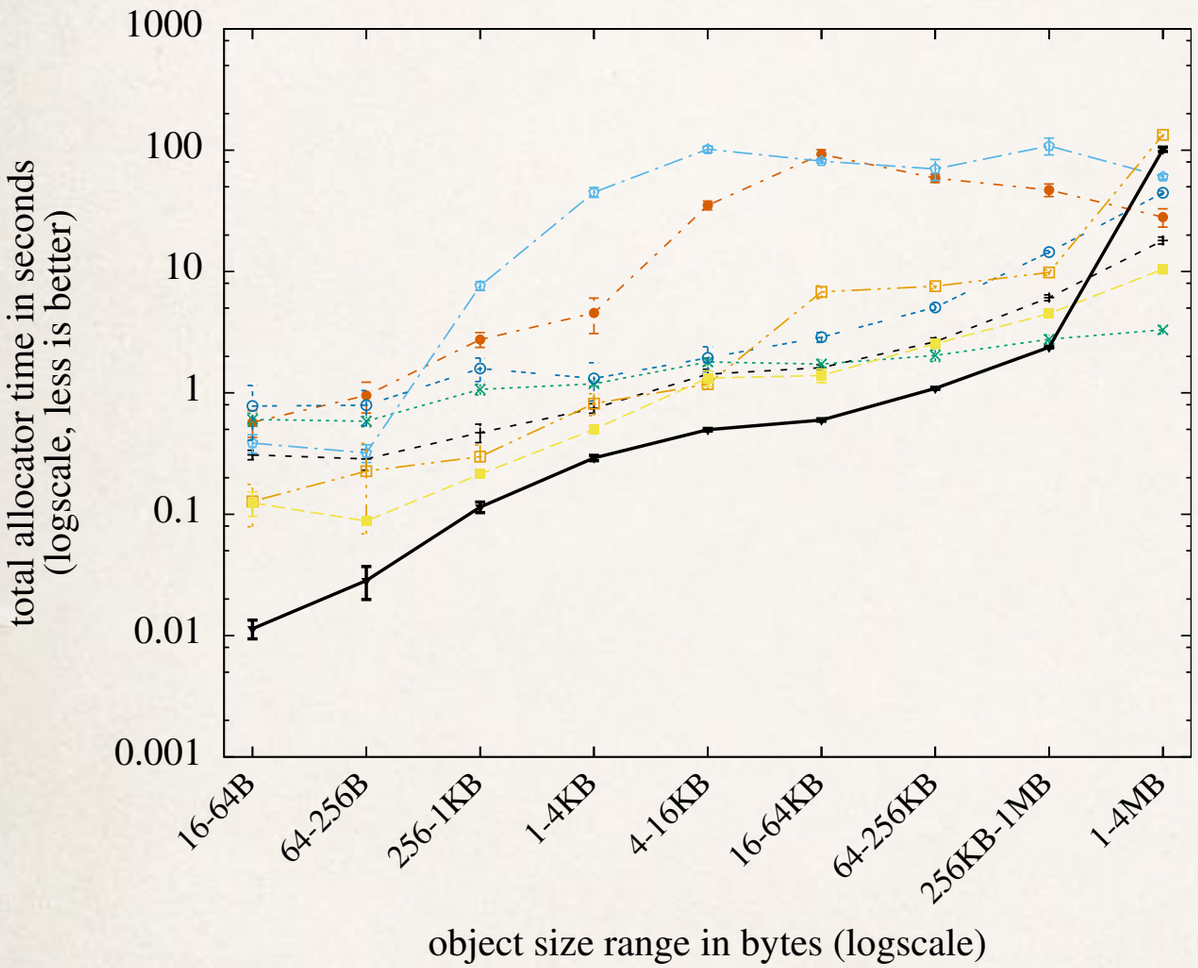
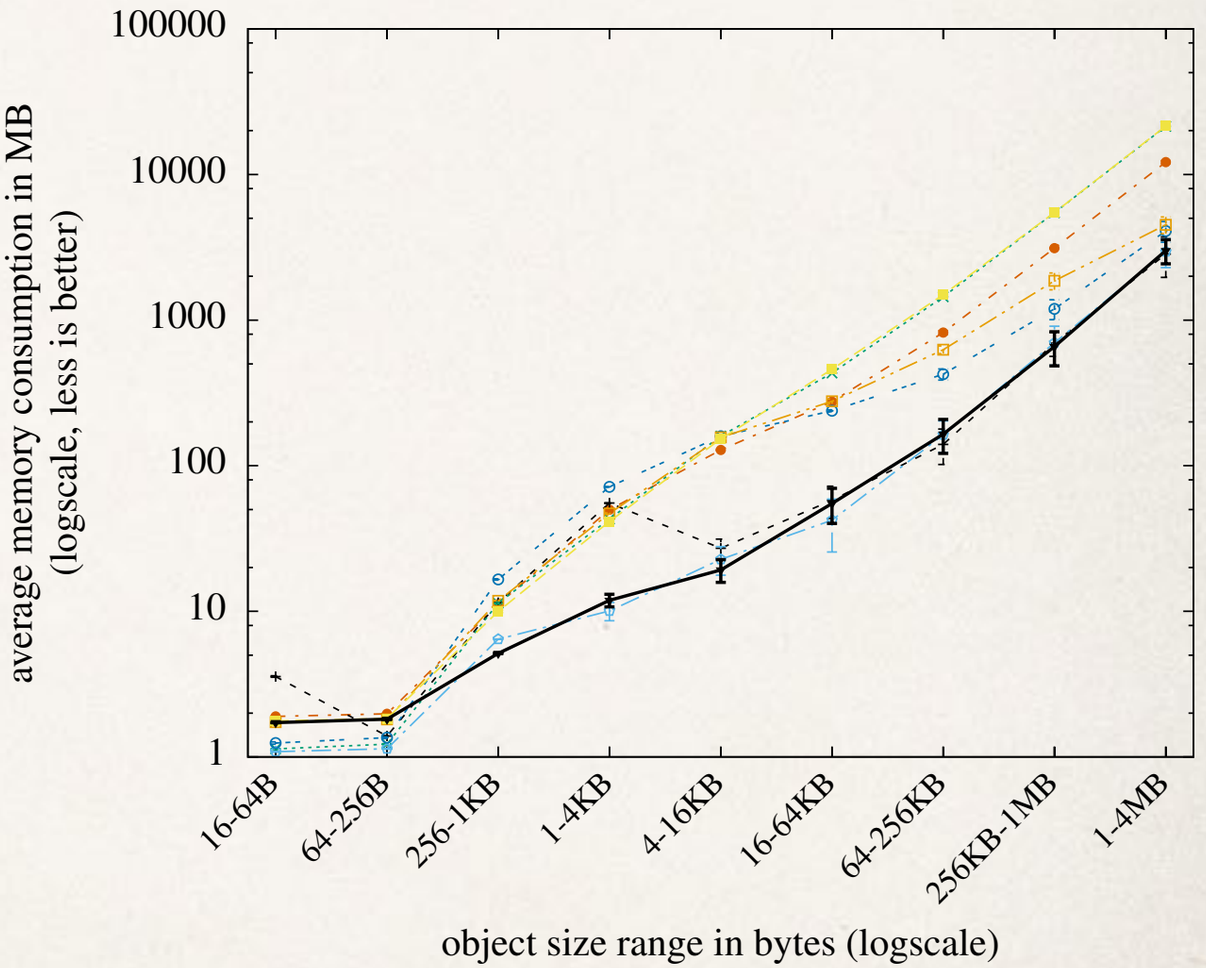


Figure 1: Structure of arena, virtual spans, and real spans

Object Size



(a) Total allocator time



(b) Average memory consumption

Figure 10: Temporal and spatial performance for the object-size robustness experiment at 40 threads

ACDC: Explorative Benchmarking Memory Management acdc.cs.uni-salzburg.at [ISMM13,DLS14]

- configurable multicore-scalable mutator for mimicking virtually any allocation, deallocation, sharing, and access behavior
- written in C, tracks with minimal overhead:

1. memory allocation time
2. memory deallocation time
3. memory consumption
4. memory access time

Memory Access

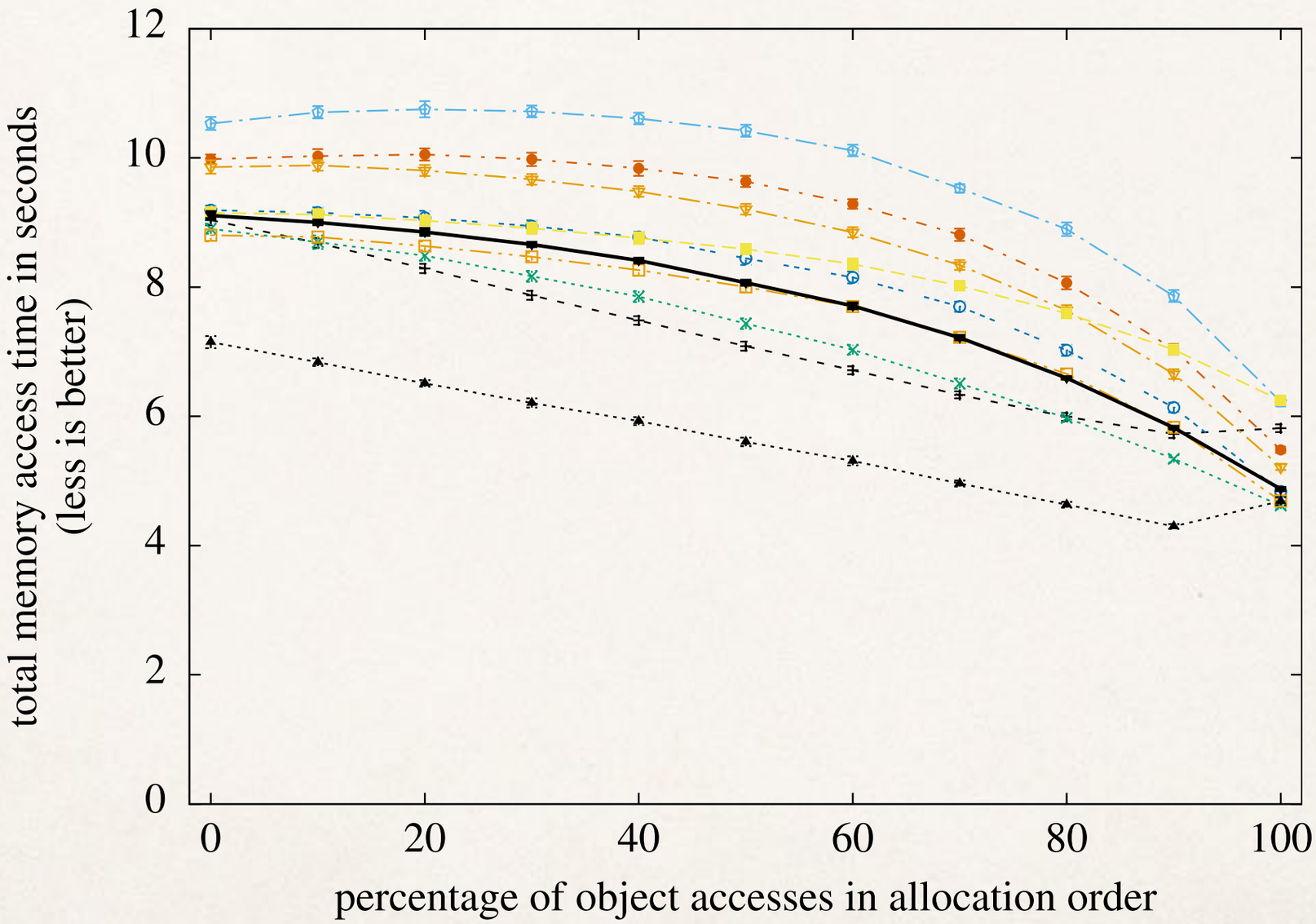


Figure 11: Memory access time for the locality experiment

How do we teach computer science to students <u>not necessarily majoring in computer science</u> but who anyway <u>code every day</u>?

–The Computer Science Education Challenge

Selfie: Teaching Computer Science [selfie.cs.uni-salzburg.at]

- *Selfie* is a self-referential 7k-line C implementation (in a single file) of:

1. a self-compiling compiler called *starc* that compiles a tiny subset of C called C Star (C*) to a tiny subset of MIPS32 called MIPSter,

2. a self-executing emulator called *mipster* that executes MIPSter code including itself when compiled with starc,

3. a self-hosting hypervisor called *hypster* that virtualizes mipster and can host all of selfie including itself,

4. a tiny C* library called *libcstar* utilized by all of selfie, and

5. a tiny, experimental SAT solver called *babysat*.

5 statements:
assignment
while
if
return
procedure()

```
int atoi(int *s) {
        int i;
        int n;
        int c;

        i = 0;
        n = 0;
        c = *(s+i);

        while (c != 0) {
            n = n * 10 + c - '0';
            if (n < 0)
                return -1;

            i = i + 1;
            c = *(s+i);
        }

        return n;
}
```

no data types other than `int` and `int*` and dereferencing: the * operator

character literals
string literals

integer arithmetics
pointer arithmetics

no bitwise operators
no Boolean operators

library: `exit`, `malloc`, `open`, `read`, `write`

```
> make
cc -w -m32 -D'main(a,b)=main(a,char**argv)' selfie.c -o selfie
```

bootstrapping `selfie.c` *into x86* `selfie` *executable*
using standard C compiler

(now also available for RISC-V machines)

```
> ./selfie
./selfie: usage: selfie { -c { source } | -o binary | -s assembly
| -l binary } [ ( -m | -d | -y | -min | -mob ) size ... ]
```

selfie usage

```
> ./selfie -c selfie.c

./selfie: this is selfie's starc compiling selfie.c

./selfie: 176408 characters read in 7083 lines and 969 comments
./selfie: with 97779(55.55%) characters in 28914 actual symbols
./selfie: 261 global variables, 289 procedures, 450 string literals
./selfie: 1958 calls, 723 assignments, 57 while, 572 if, 243 return
./selfie: 121660 bytes generated with 28779 instructions and 6544
bytes of data
```

compiling `selfie.c` *with x86* `selfie` *executable*

(takes seconds)

```
> ./selfie -c selfie.c -m 2 -c selfie.c

./selfie: this is selfie's starc compiling selfie.c

./selfie: this is selfie's mipster executing selfie.c with 2MB of
physical memory

selfie.c: this is selfie's starc compiling selfie.c

selfie.c: exiting with exit code 0 and 1.05MB of mallocated memory

./selfie: this is selfie's mipster terminating selfie.c with exit code
0 and 1.16MB of mapped memory
```

compiling `selfie.c` *with x86* `selfie` *executable into a MIPSter executable*

and

then running that MIPSter executable to compile `selfie.c` *again*

(takes ~6 minutes)

```
> ./selfie -c selfie.c -o selfie1.m -m 2 -c selfie.c -o selfie2.m

./selfie: this is selfie's starc compiling selfie.c
./selfie: 121660 bytes with 28779 instructions and 6544 bytes of data
written into selfie1.m

./selfie: this is selfie's mipster executing selfie1.m with 2MB of
physical memory

selfie1.m: this is selfie's starc compiling selfie.c
selfie1.m: 121660 bytes with 28779 instructions and 6544 bytes of data
written into selfie2.m

selfie1.m: exiting with exit code 0 and 1.05MB of mallocated memory

./selfie: this is selfie's mipster terminating selfie1.m with exit
code 0 and 1.16MB of mapped memory
```

compiling `selfie.c` *into a MIPSter executable* `selfie1.m`

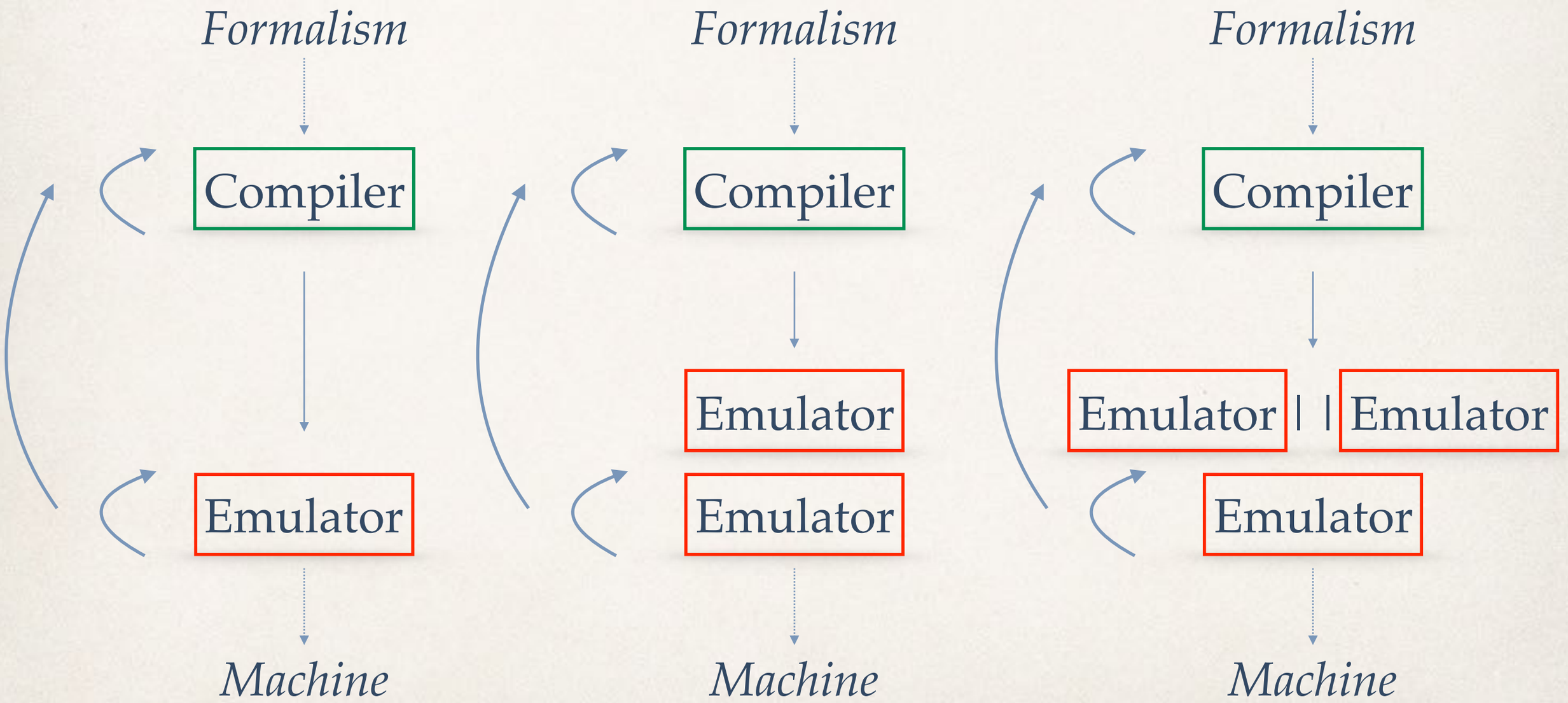
and

then running `selfie1.m` *to compile* `selfie.c`

into another MIPSter executable `selfie2.m`

(takes ~6 minutes)

Sandboxed Concurrency: 1-Week Homework Assignment



```
> ./selfie -c selfie.c -m 2 -c selfie.c -m 2 -c selfie.c
```

compiling `selfie.c` *with x86* `selfie` executable

and

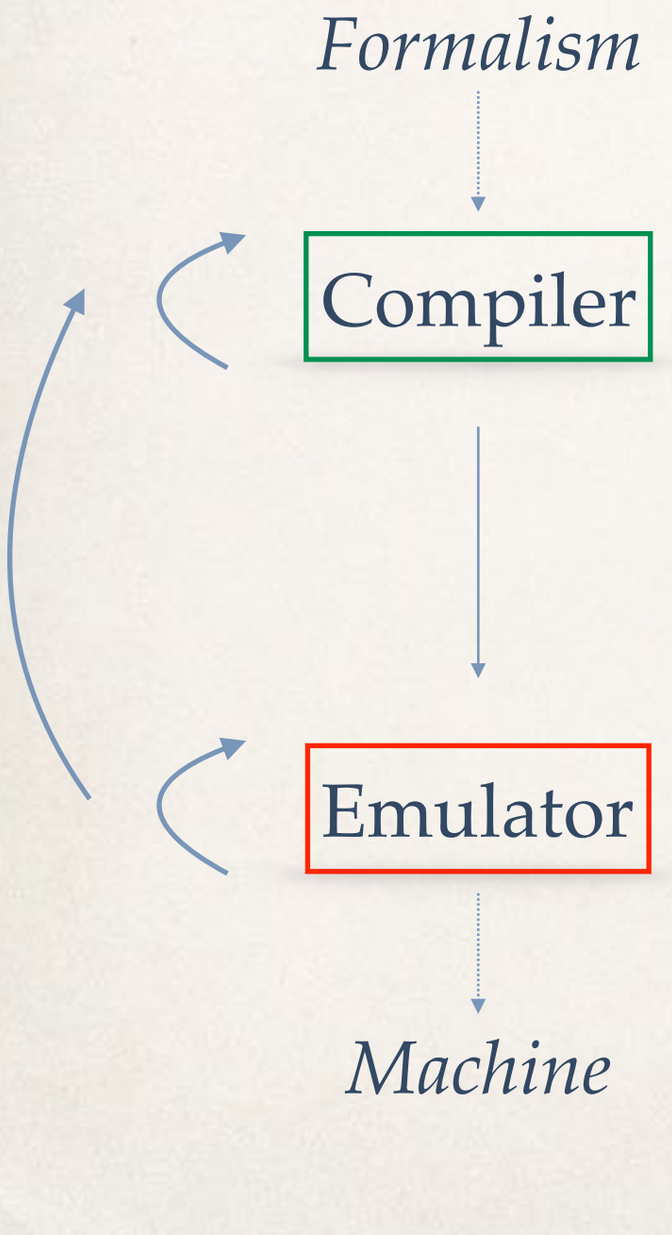
then running that executable to compile `selfie.c` *again*

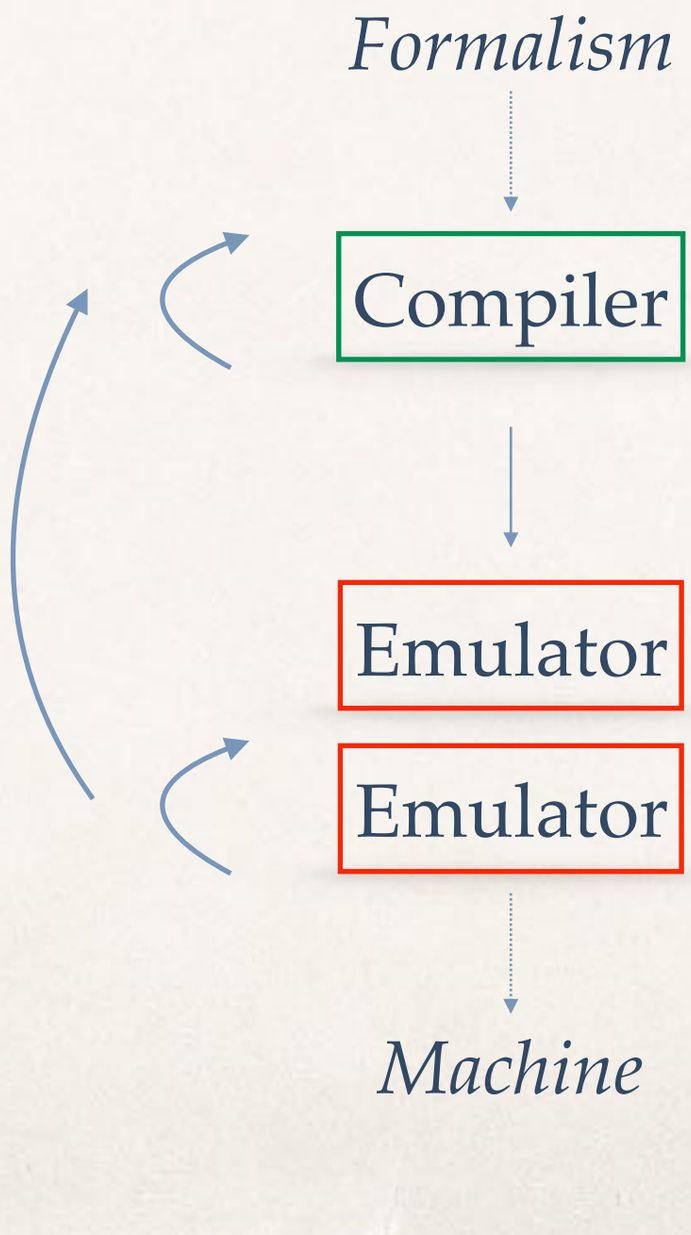
and

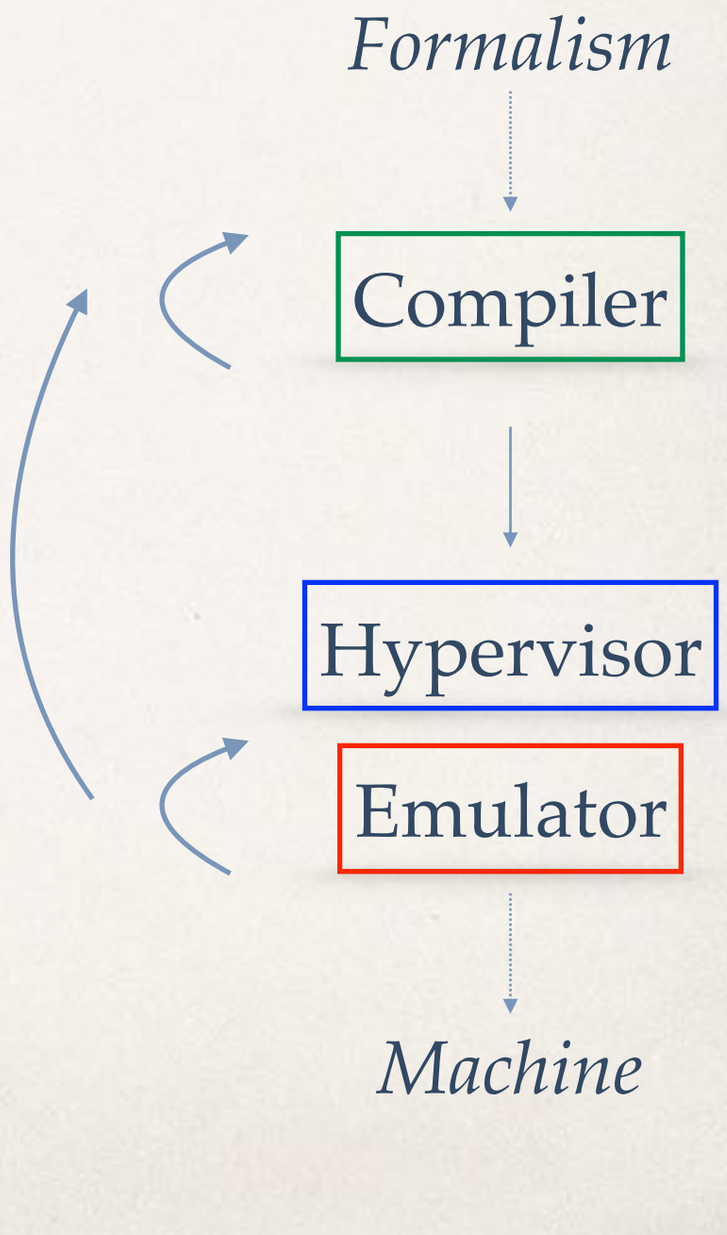
then running that executable to compile `selfie.c` *again*

(takes ~24 hours)

Emulation versus Virtualization







```
> ./selfie -c selfie.c -m 2 -c selfie.c -y 2 -c selfie.c
```

compiling `selfie.c` *with x86* `selfie` *executable*

and

then running that executable to compile `selfie.c` *again*

and

then ***hosting*** *that executable in a virtual machine to compile* `selfie.c` *again*

(takes ~12 minutes)

"How do we introduce self-model-checking and maybe even self-verification into Selfie?"

https://github.com/cksystemsteaching/selfie/tree/vipster

personal cups

Coffee Machine - Full Clean - Roster

No coffee, No workee.

GUATEMALA

Thank you!