

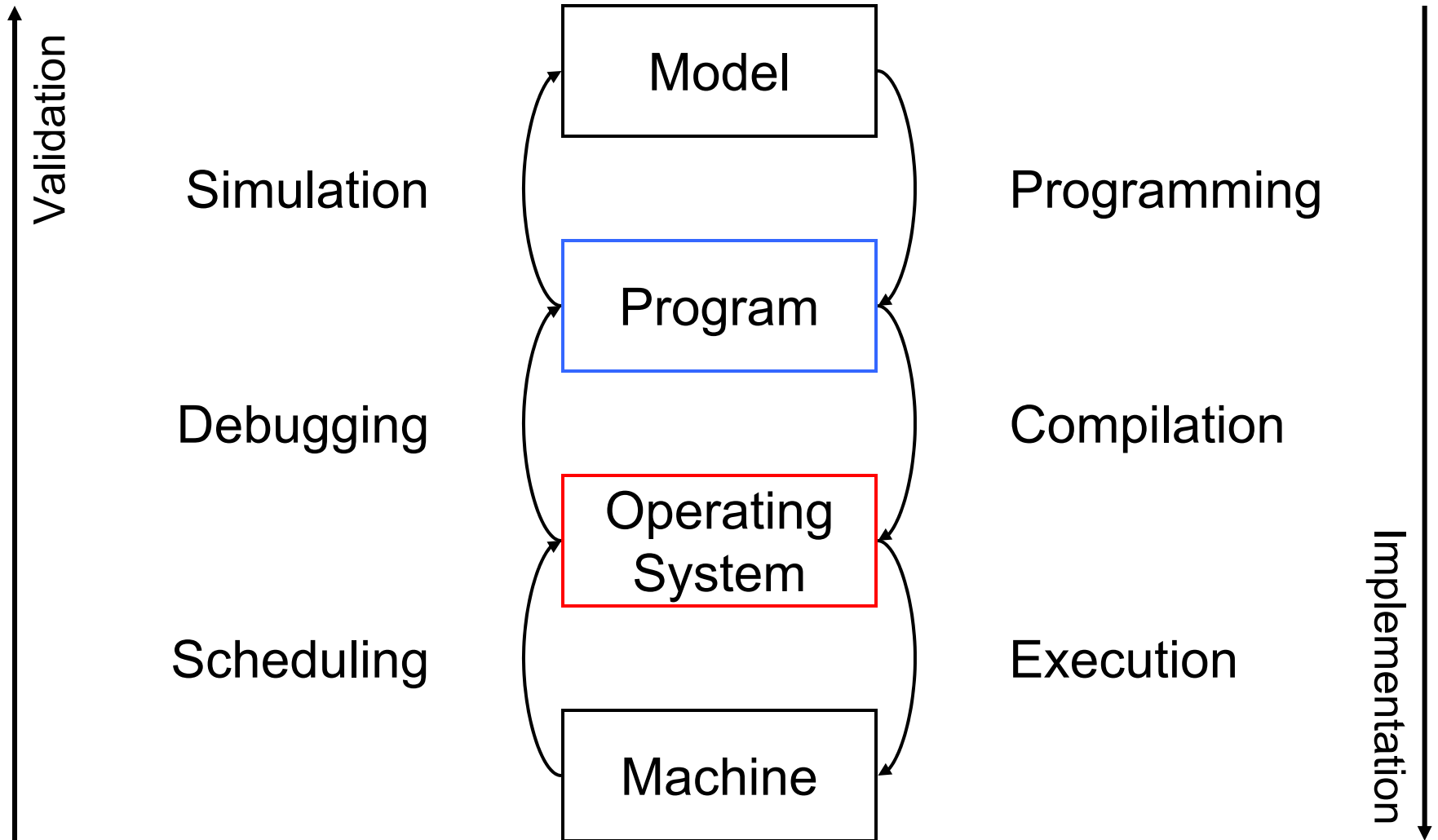
Real-Time Programming Based on Schedule-Carrying Code

Christoph M. Kirsch

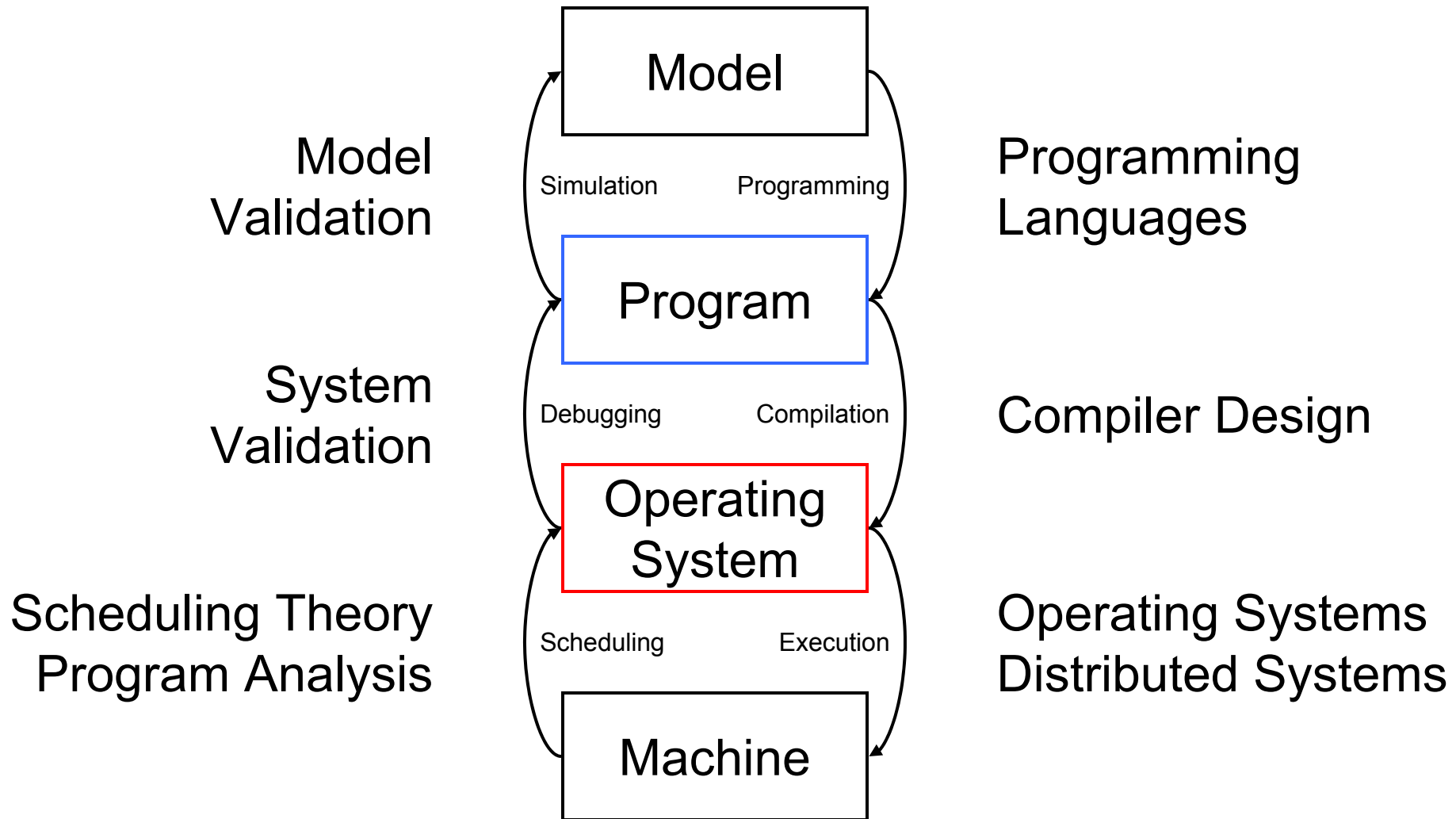
UC Berkeley

www.eecs.berkeley.edu/~cm

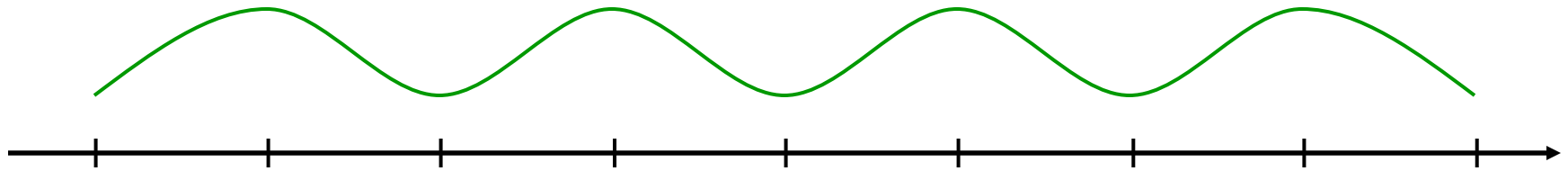
Outline



Research Areas



Real Time vs. Soft Time

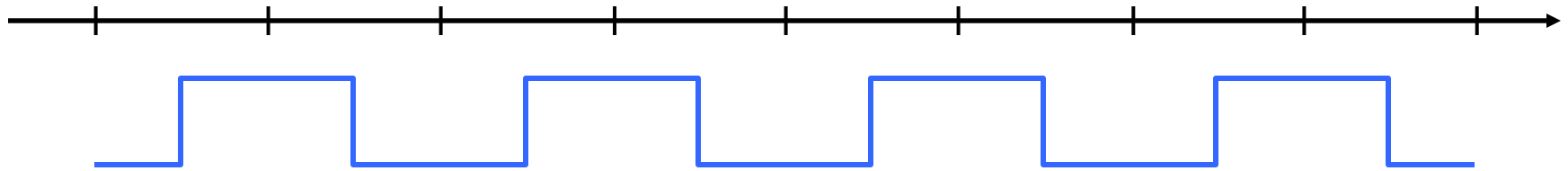


Environment Processes

Real Time



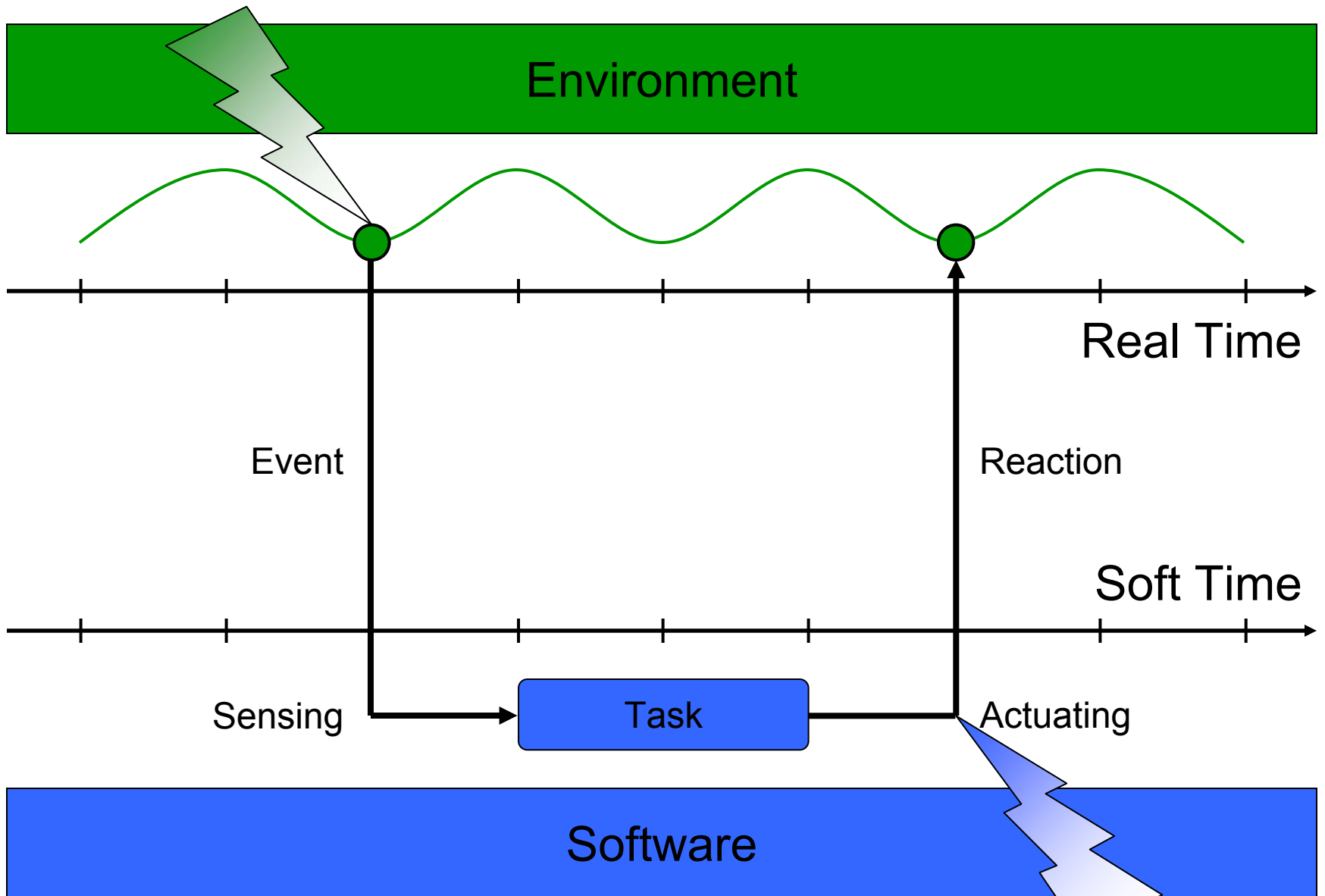
Software Processes



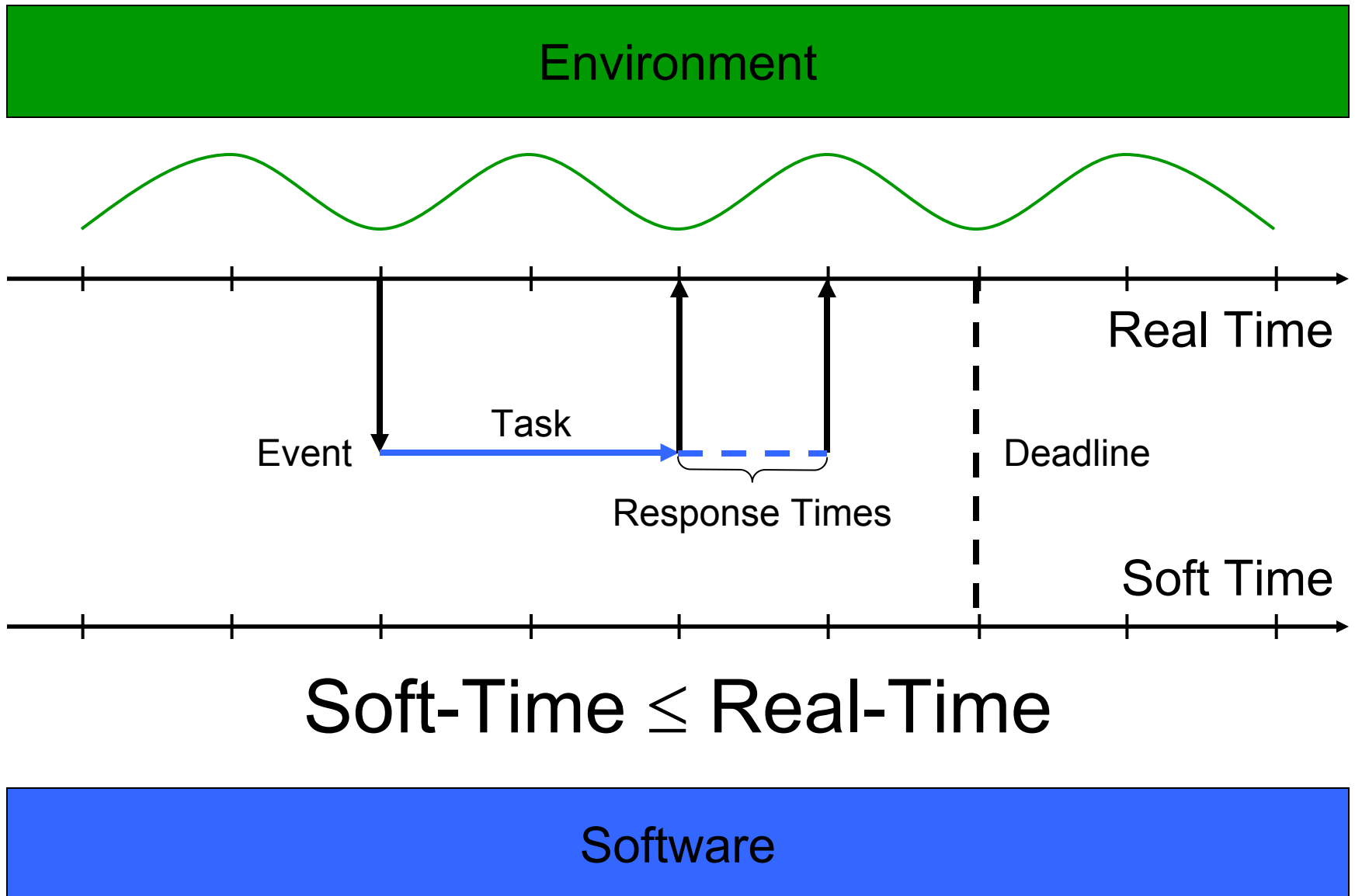
Soft Time



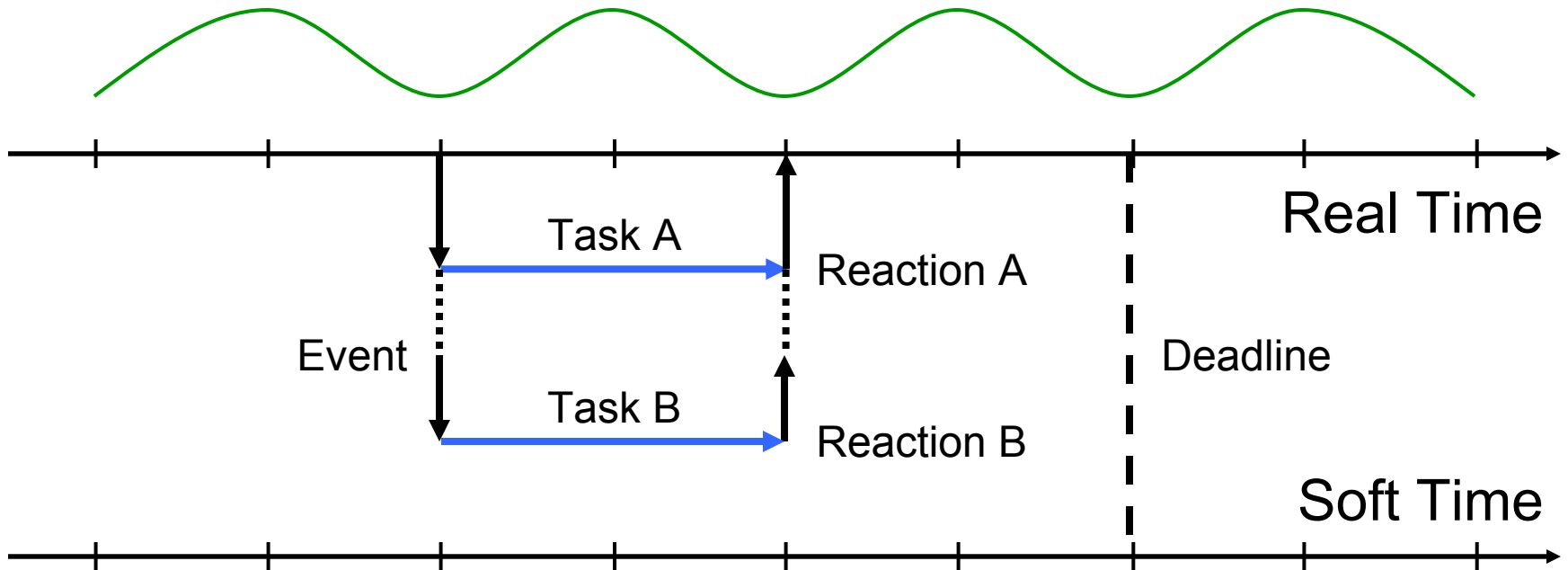
Sensing, Computing, Actuating



Real-Time Programming



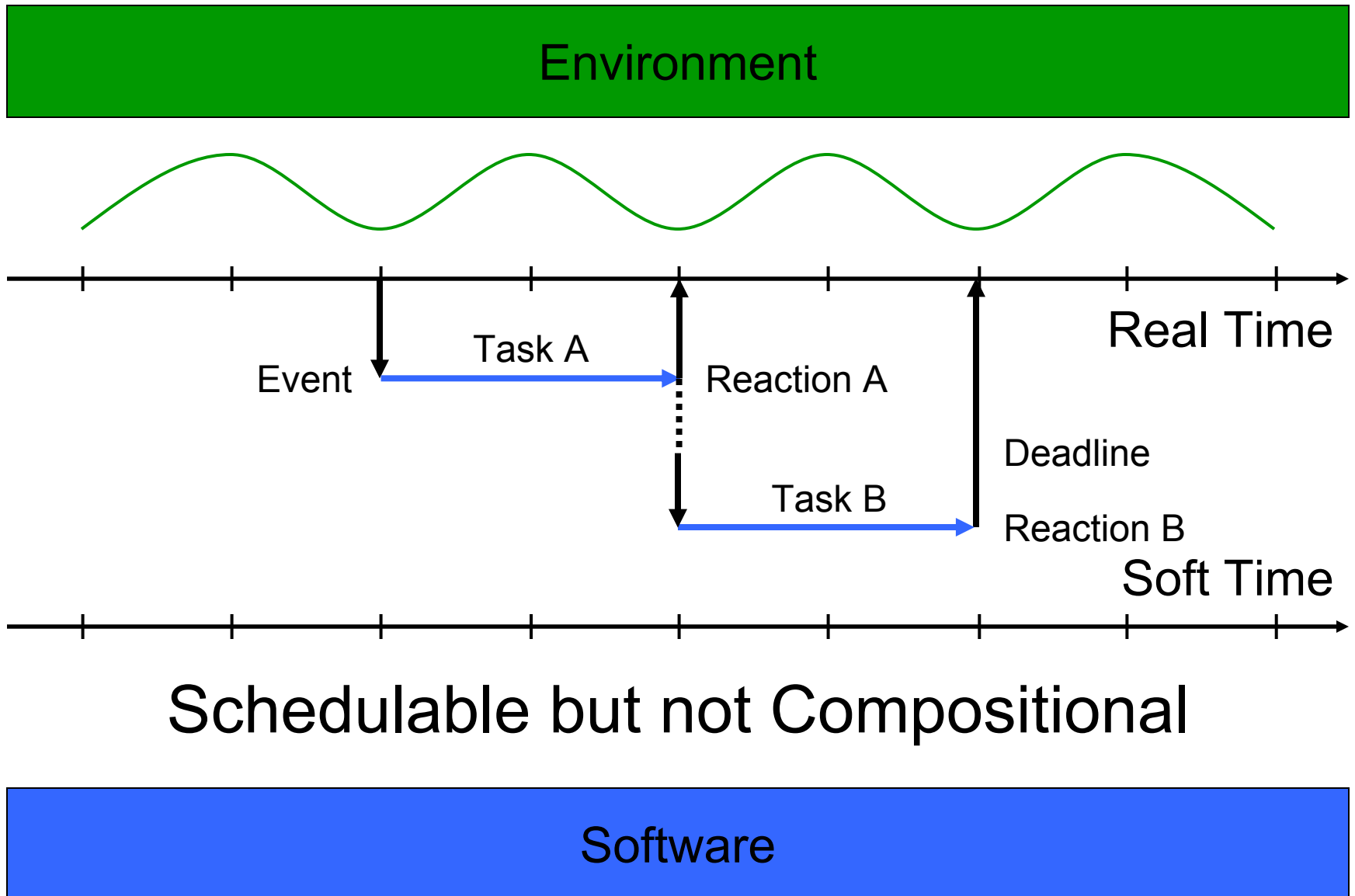
Concurrency



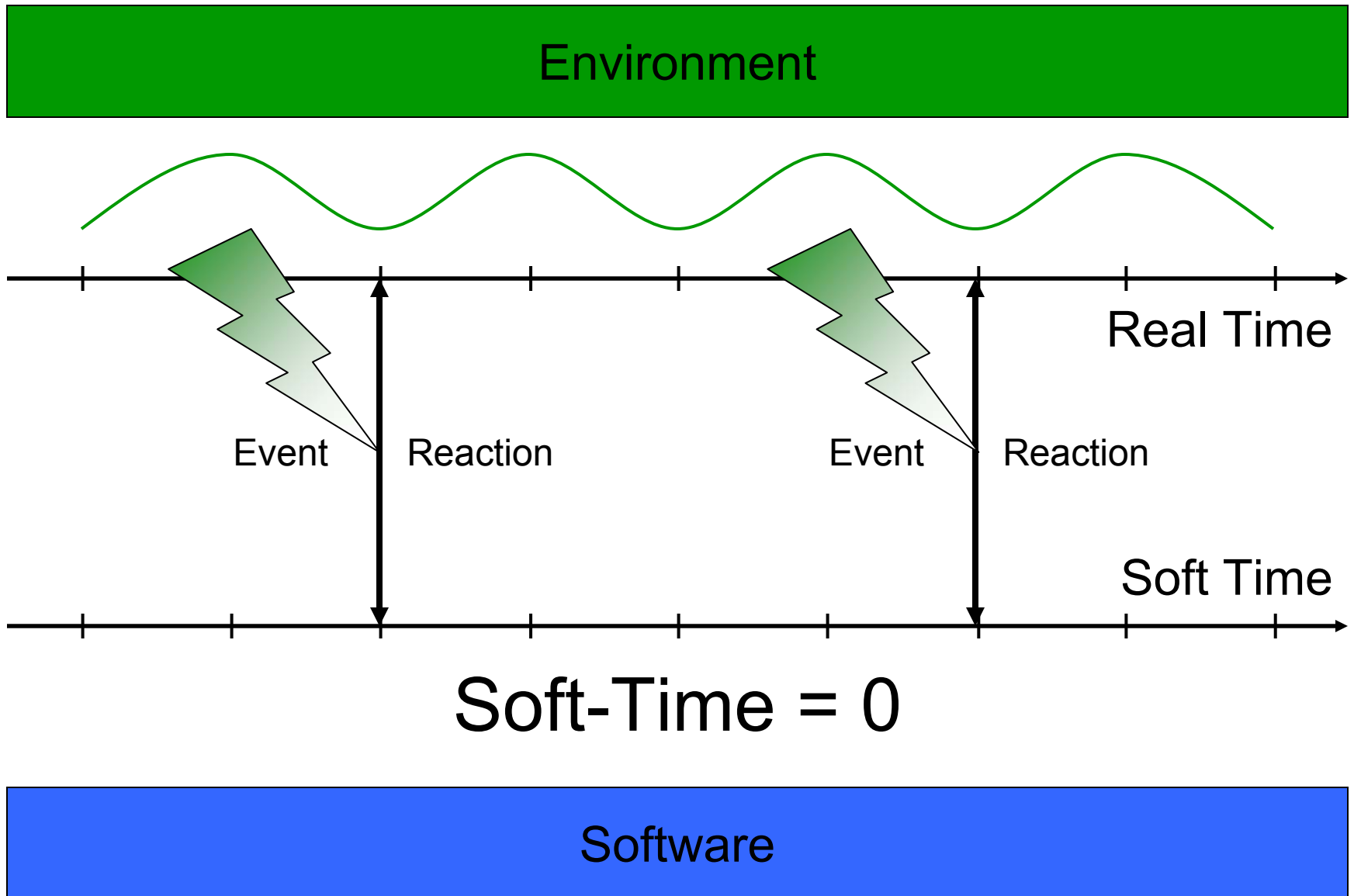
$$\text{Soft-Time} \leq \text{Real-Time}$$



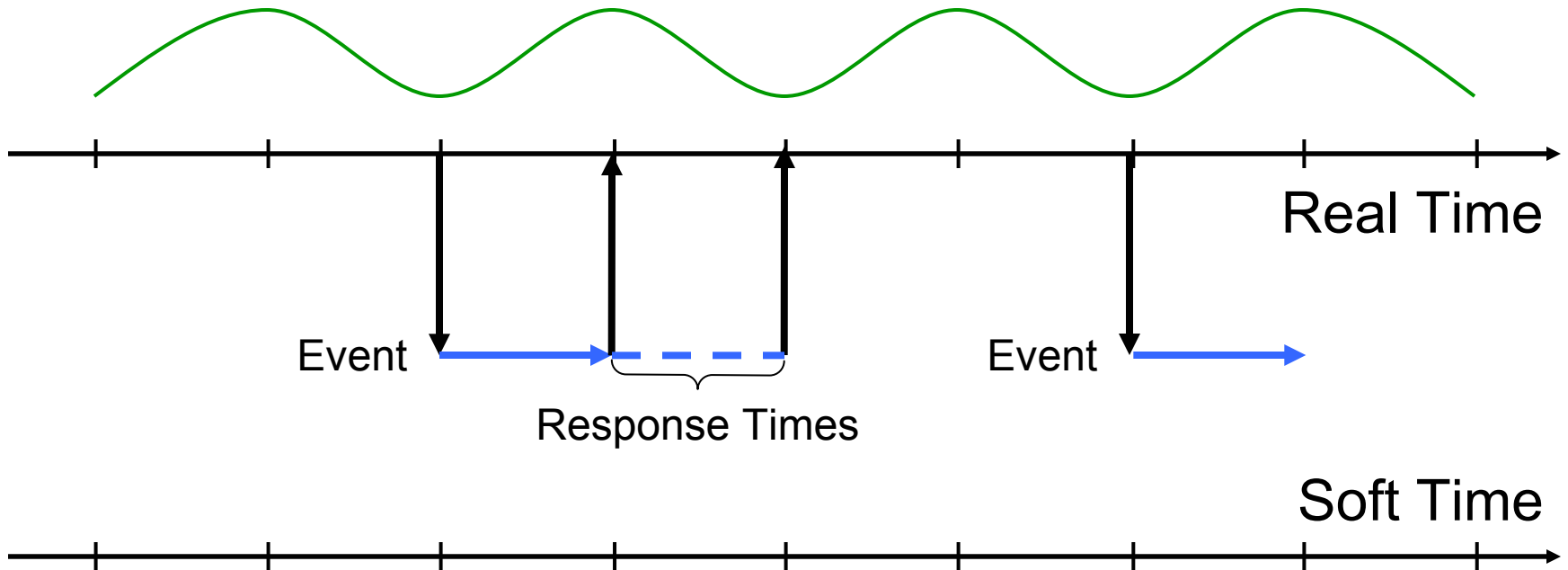
The Problem



The Synchronous Model



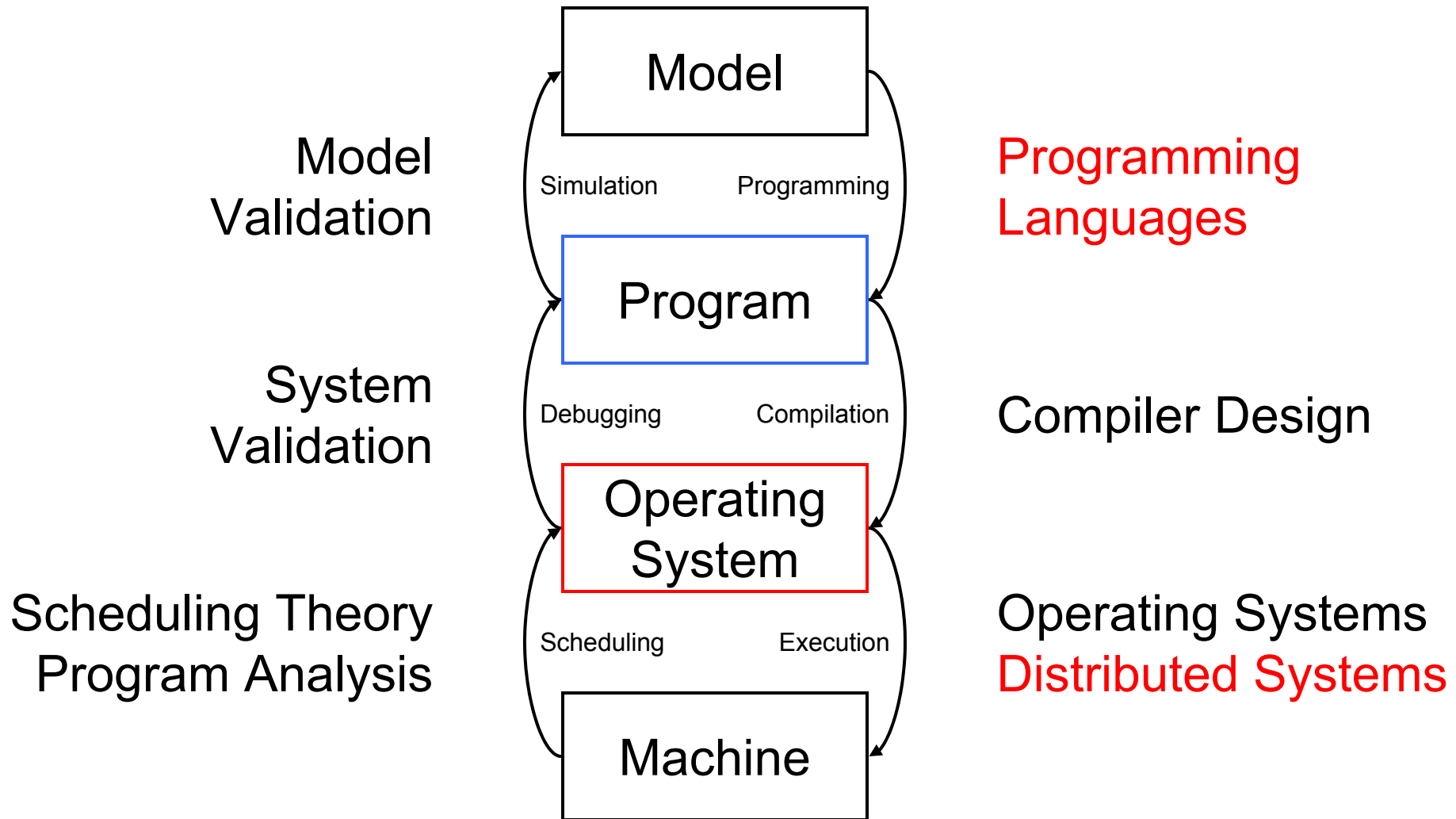
A Synchronous Implementation



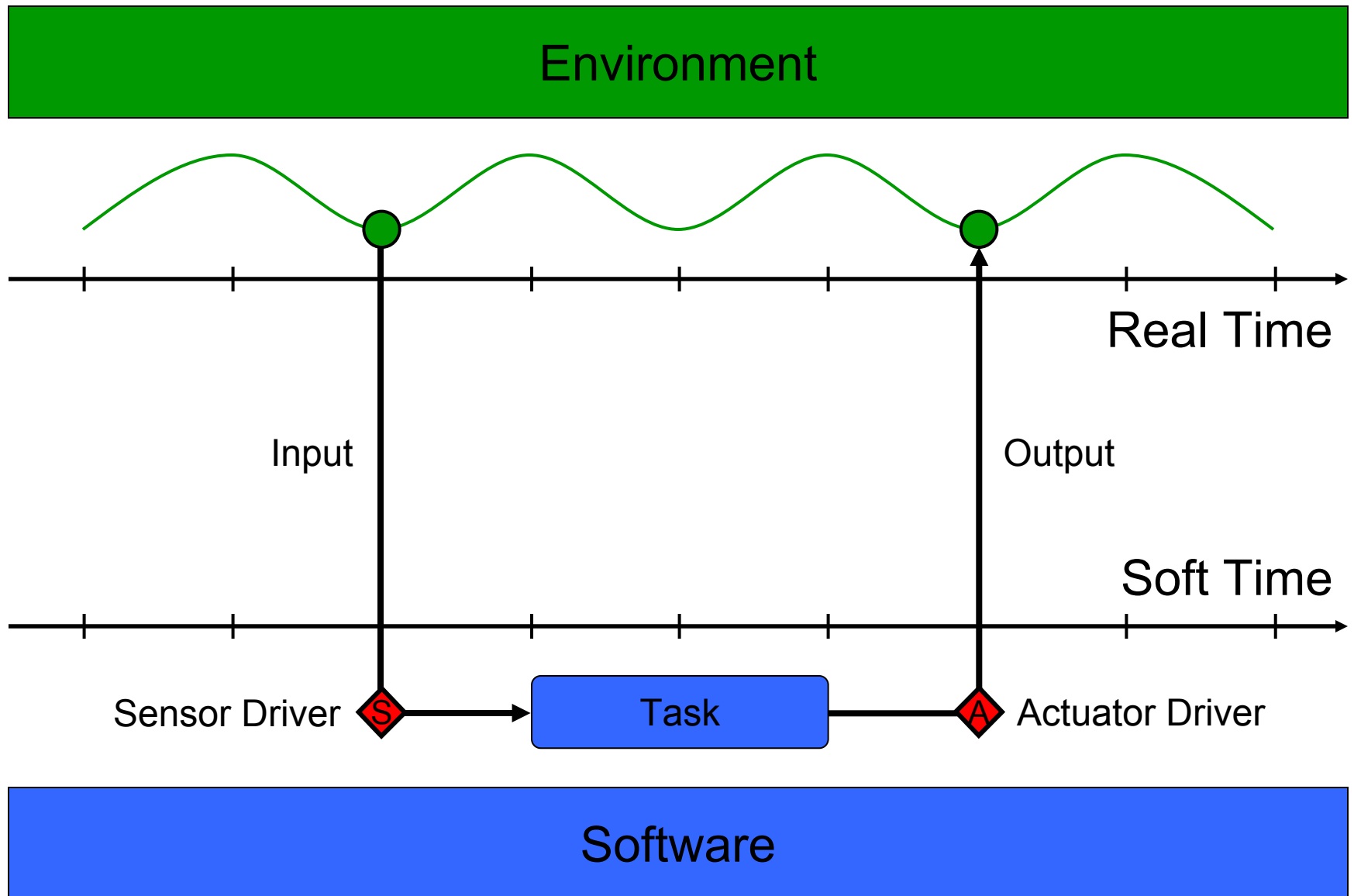
Synchronous but not Compositional



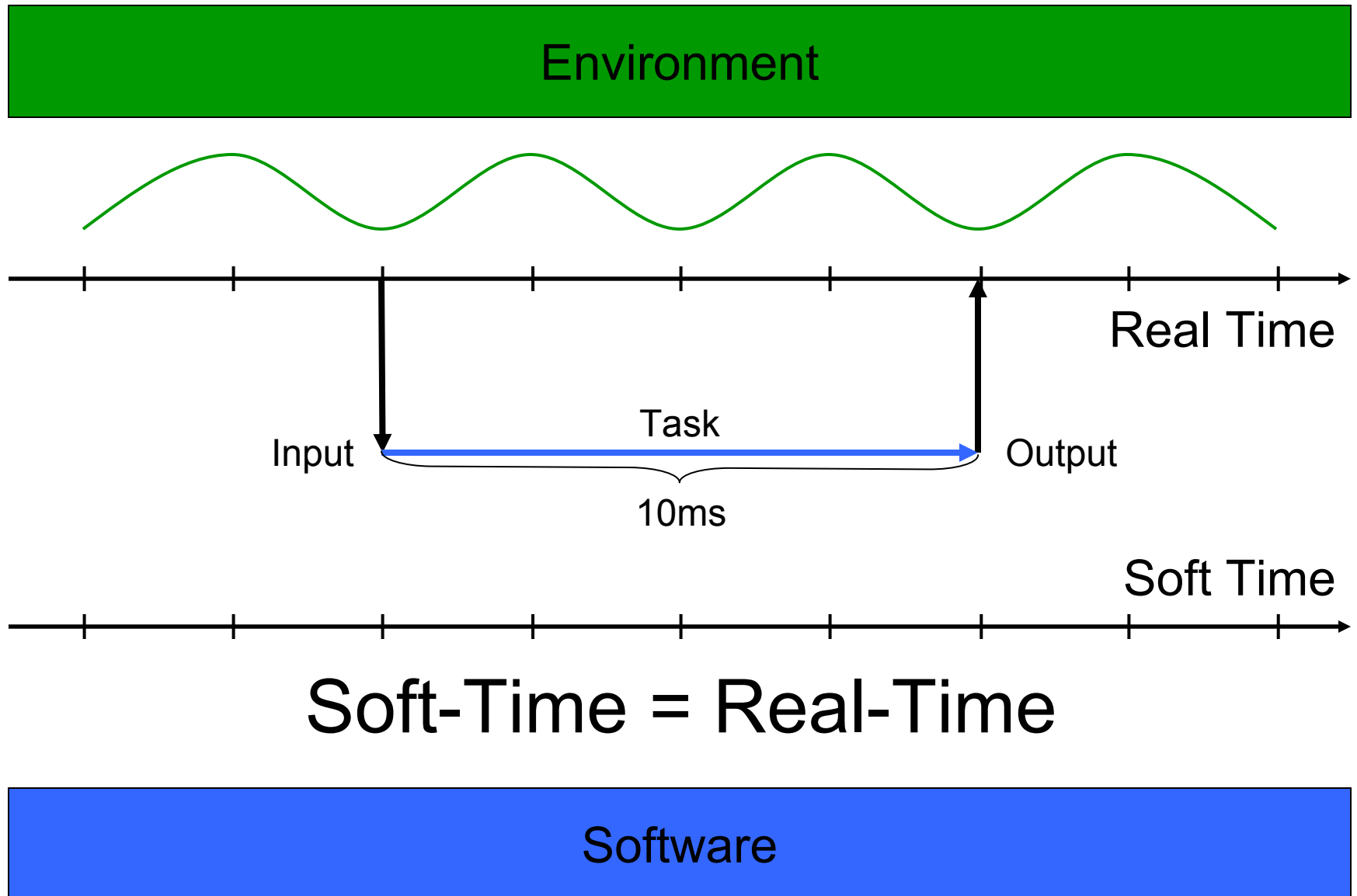
Embedded Programming Language: Giotto



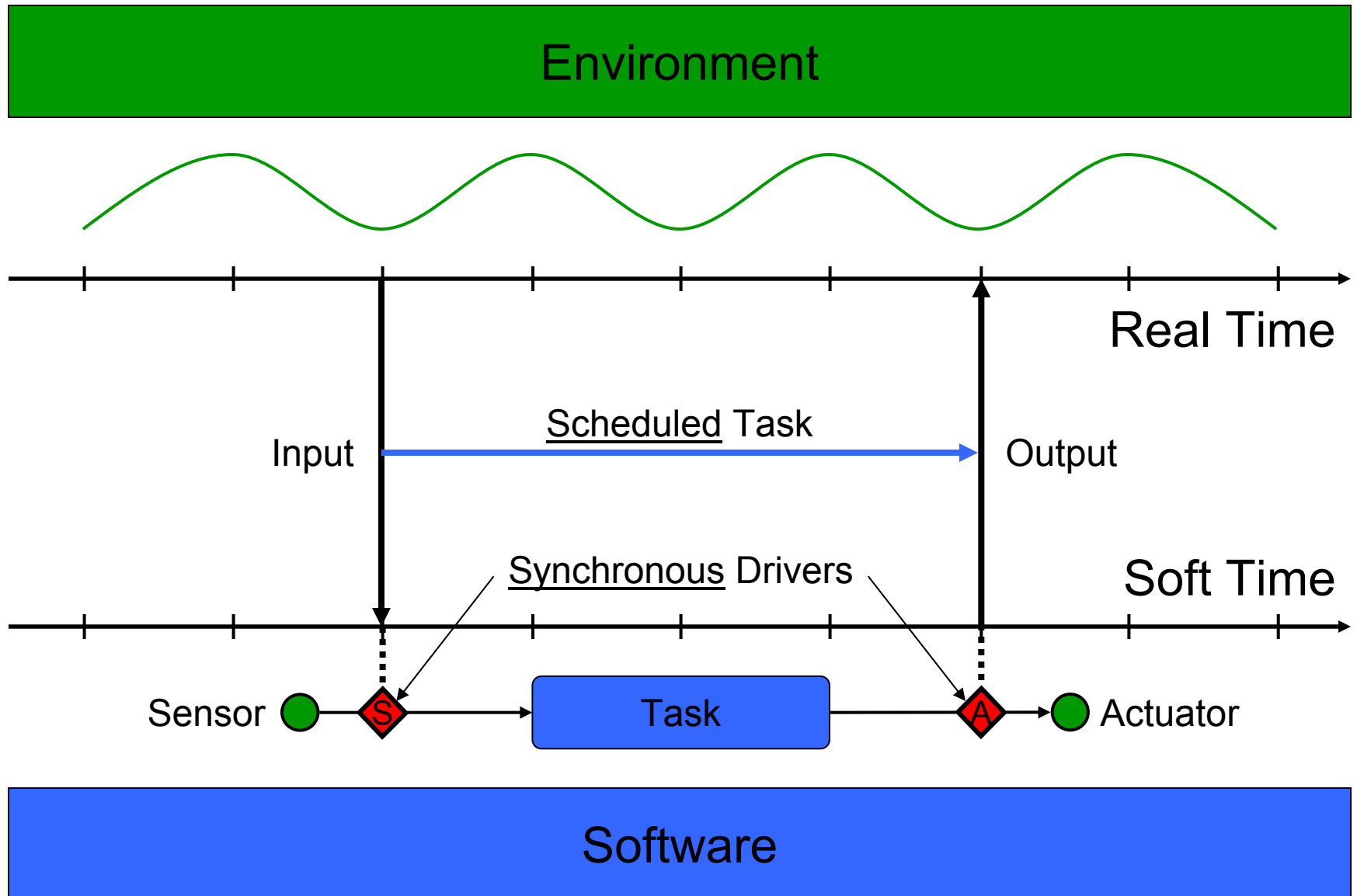
Sensing, Computing, Actuating



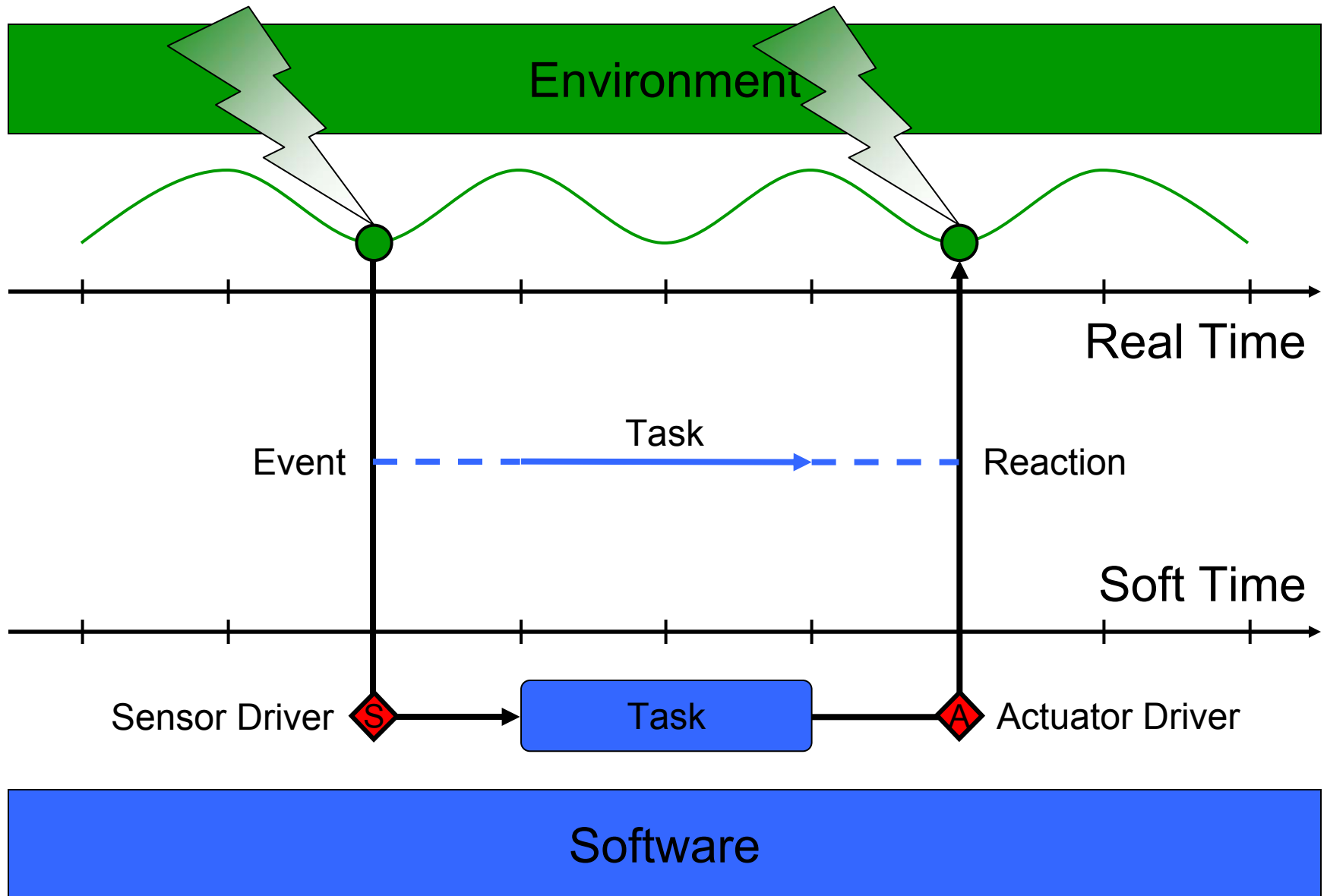
Giotto's Programming Abstraction



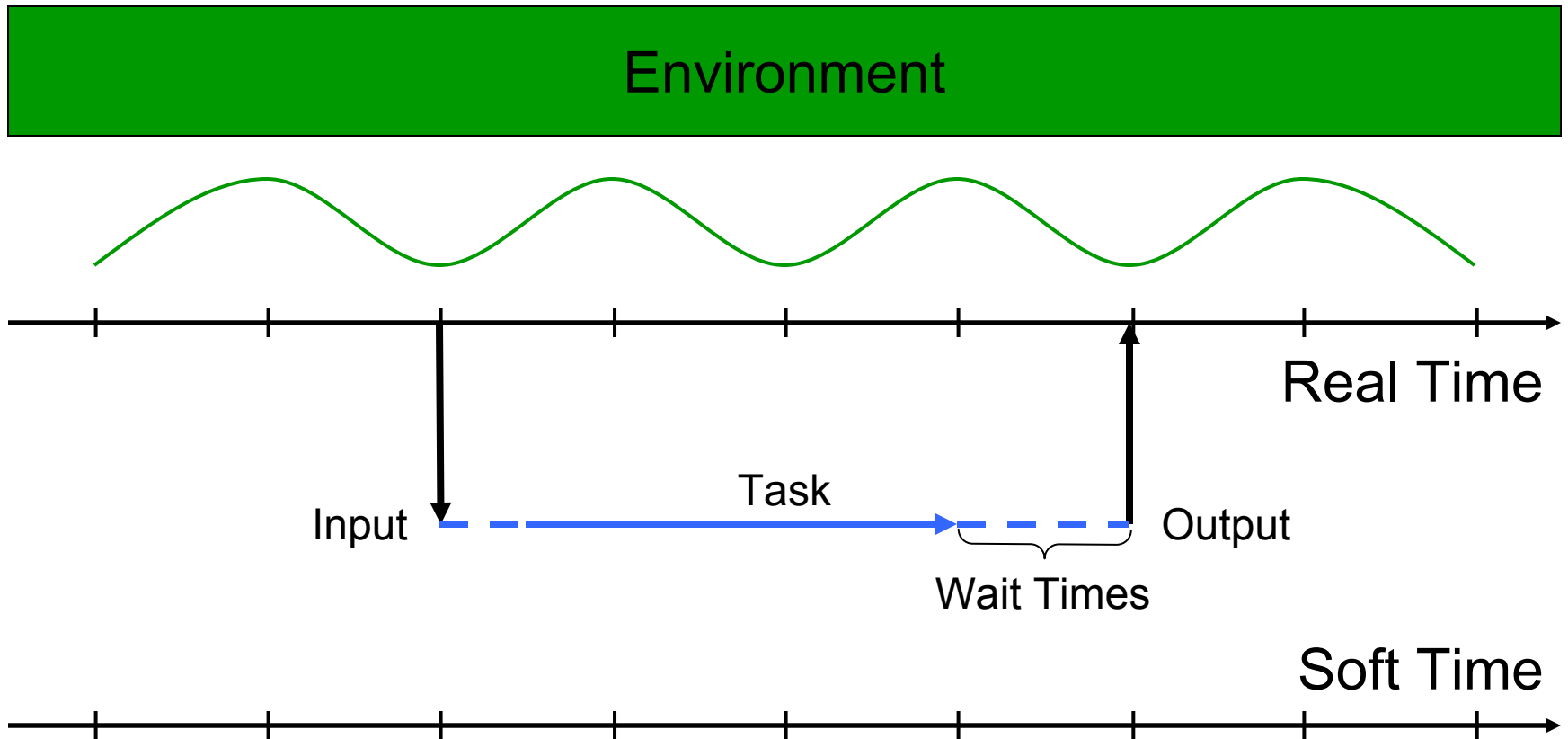
Synchronous vs. Scheduled Computation



Environment-triggered Programs



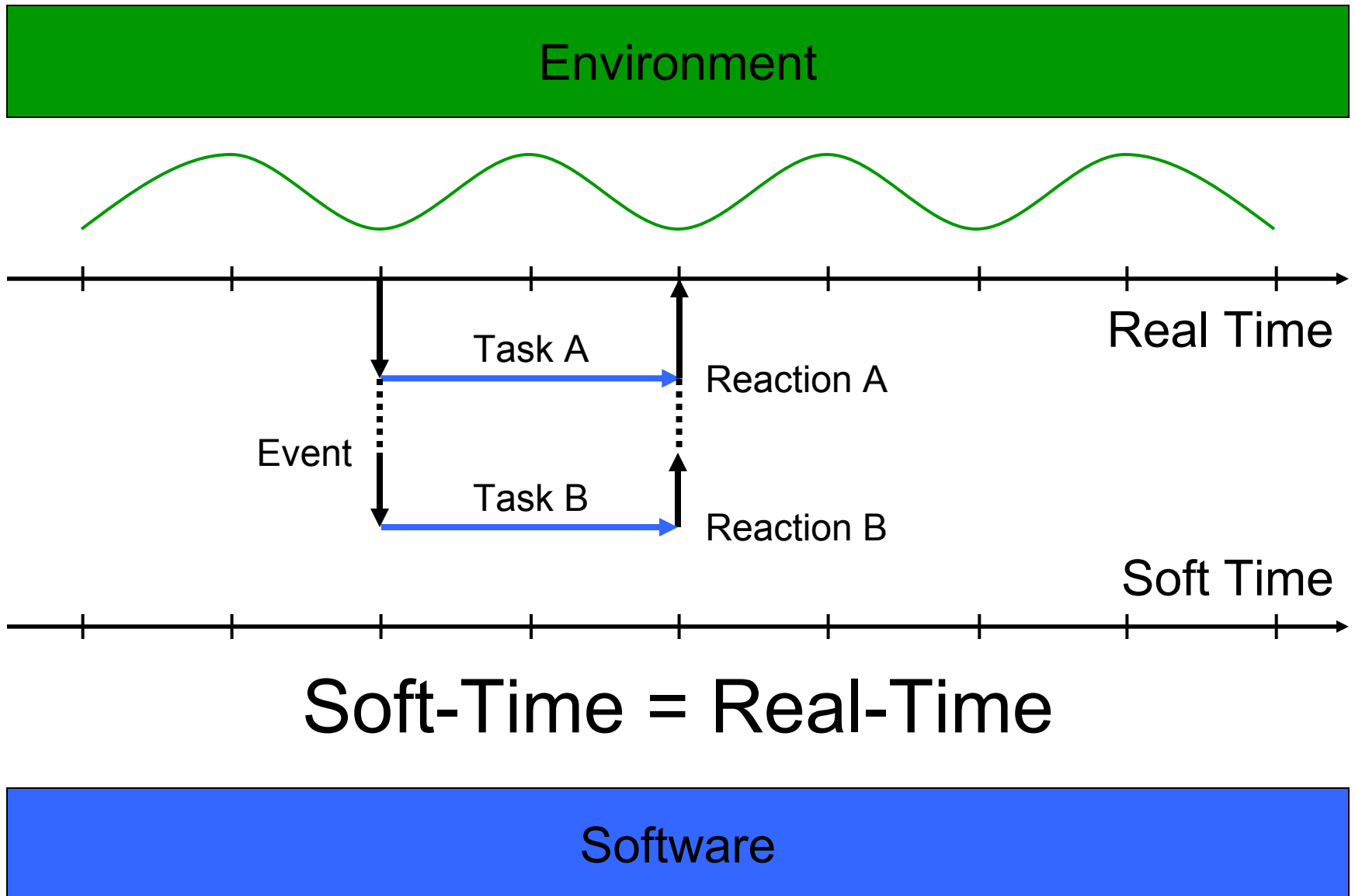
A Time-Safe Implementation



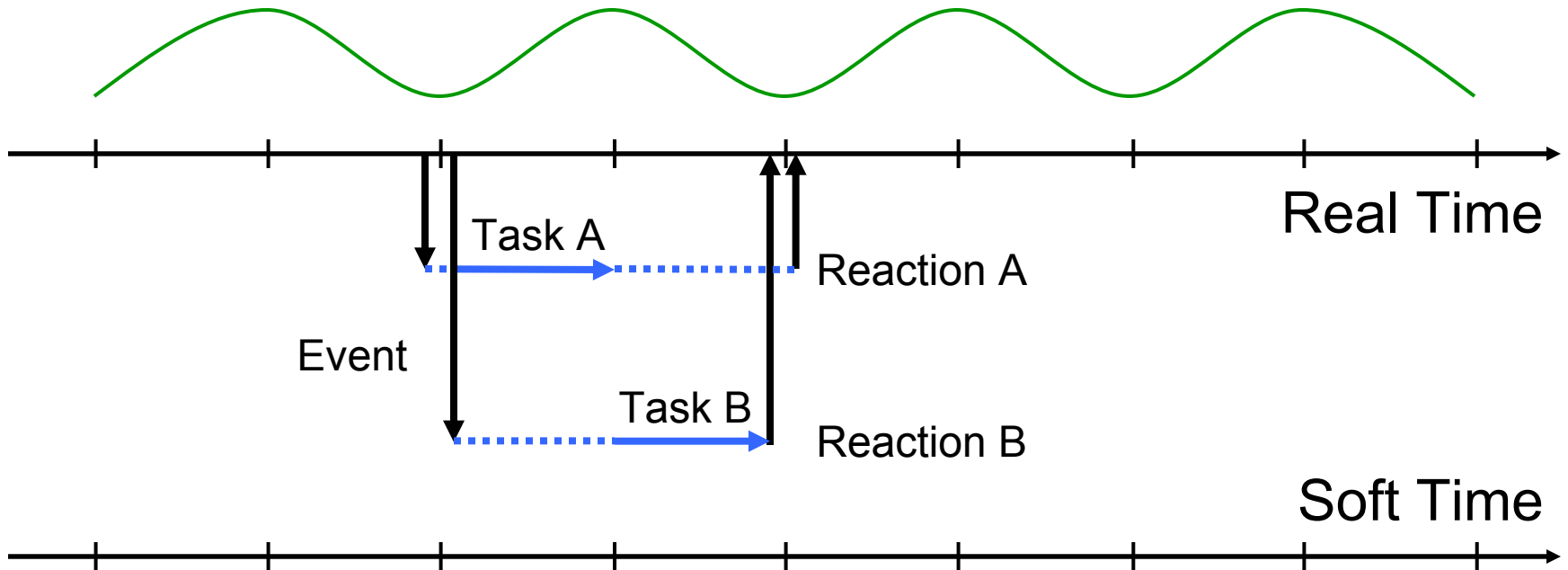
Time Safety Implies Time Determinism

Software

Abstraction



Implementation



Compositional wrt. Time Determinism



Giotto: A Time-triggered Language for Embd. Programming (Henzinger, Horowitz, Kirsch in the Proceedings of the IEEE, Jan 2003)

Giotto is a **time-triggered programming language** that supports the development of **embedded control systems**

Time-safe Giotto programs are:

- *predictable* – **deterministic** real-time code
- *platform-independent* – runs on **distributed** systems
- *multi-modal* – supports **mode switching**
- *composable* – supports **modular** compilation

Giotto is available for:

- **Linux, Windows, OSEKWorks, HelyOS**
- **Java**

Giotto on the ETH Zürich Helicopter

(Kirsch, Sanvido, Henzinger, Pree in Proc. of EMSOFT 2002)



6 degrees of freedom, 1 processor (StrongARM 200Mhz)


Giotto on the UC Berkeley Helicopter

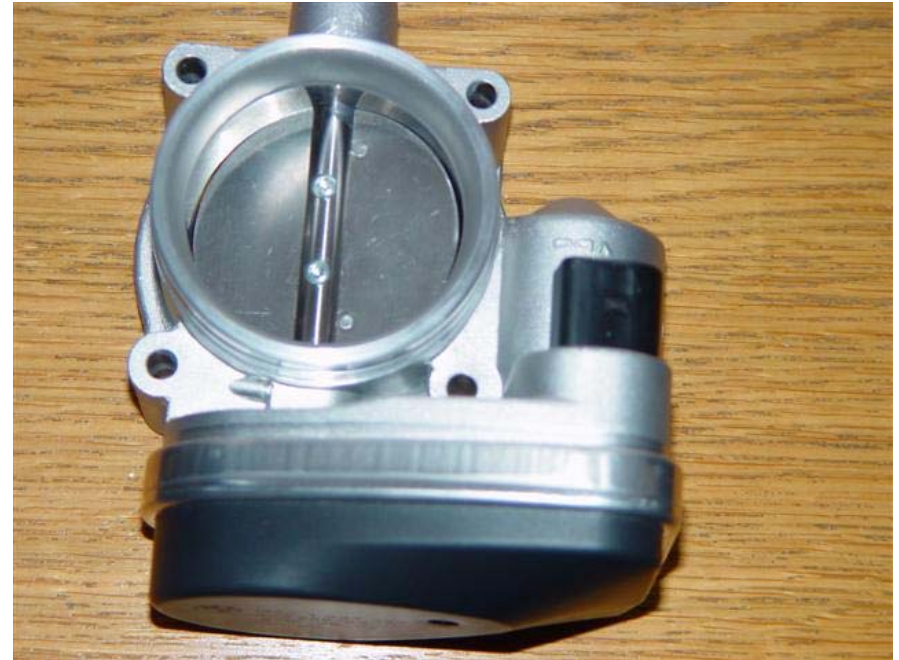
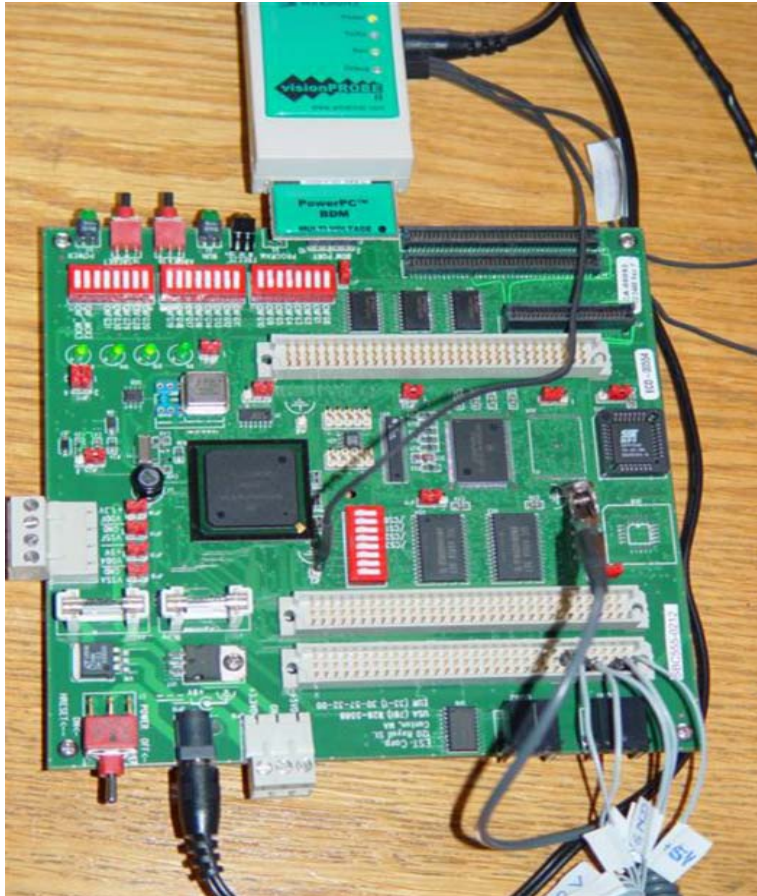
(Part of the  SEC Project with Boeing and Honeywell)



6 degrees of freedom, 3 processors (Intel x86)

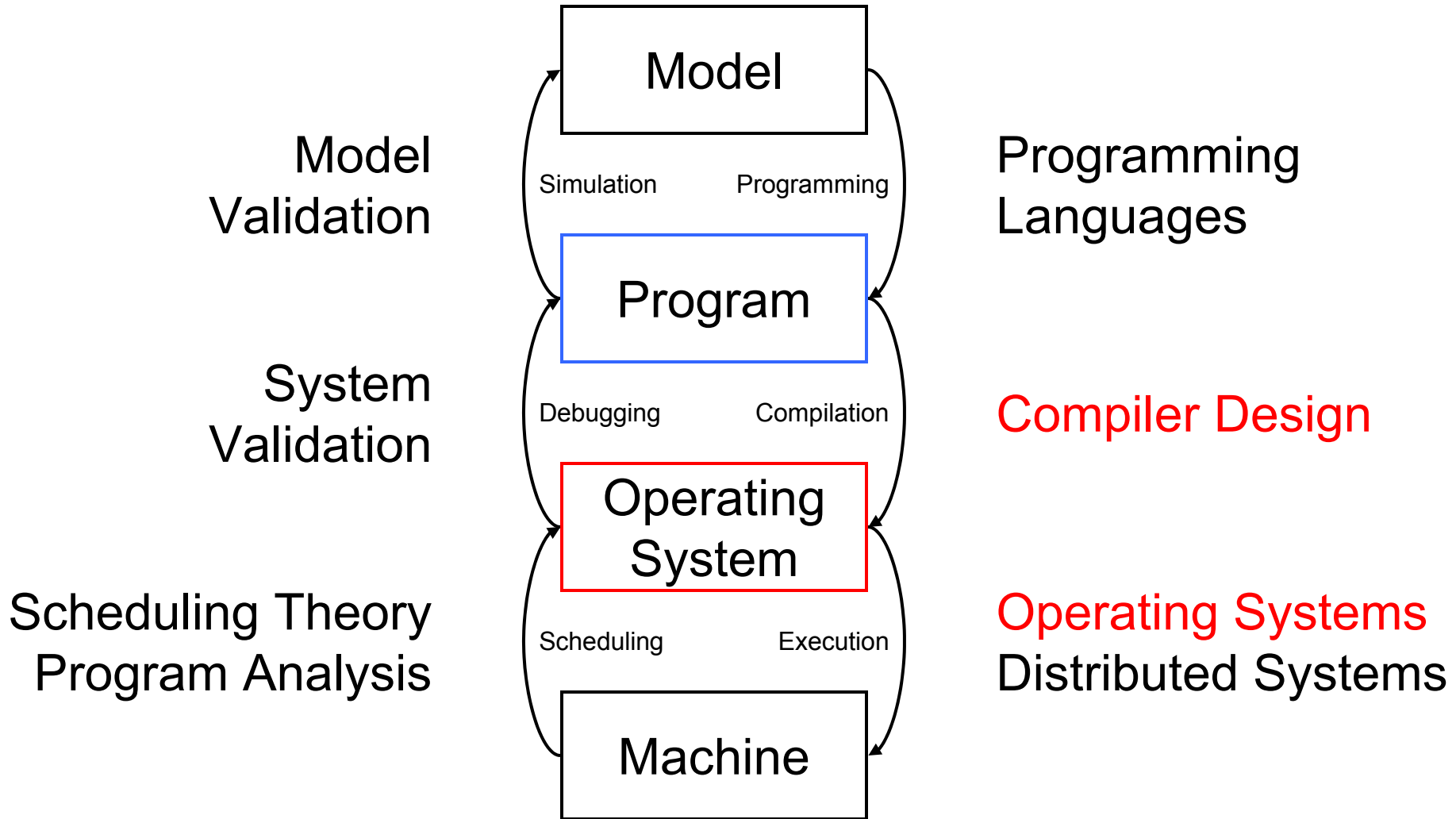
Giotto for a Drive-By-Wire BMW Throttle

(Part of the  MoBIES Project continued at Universität Salzburg)

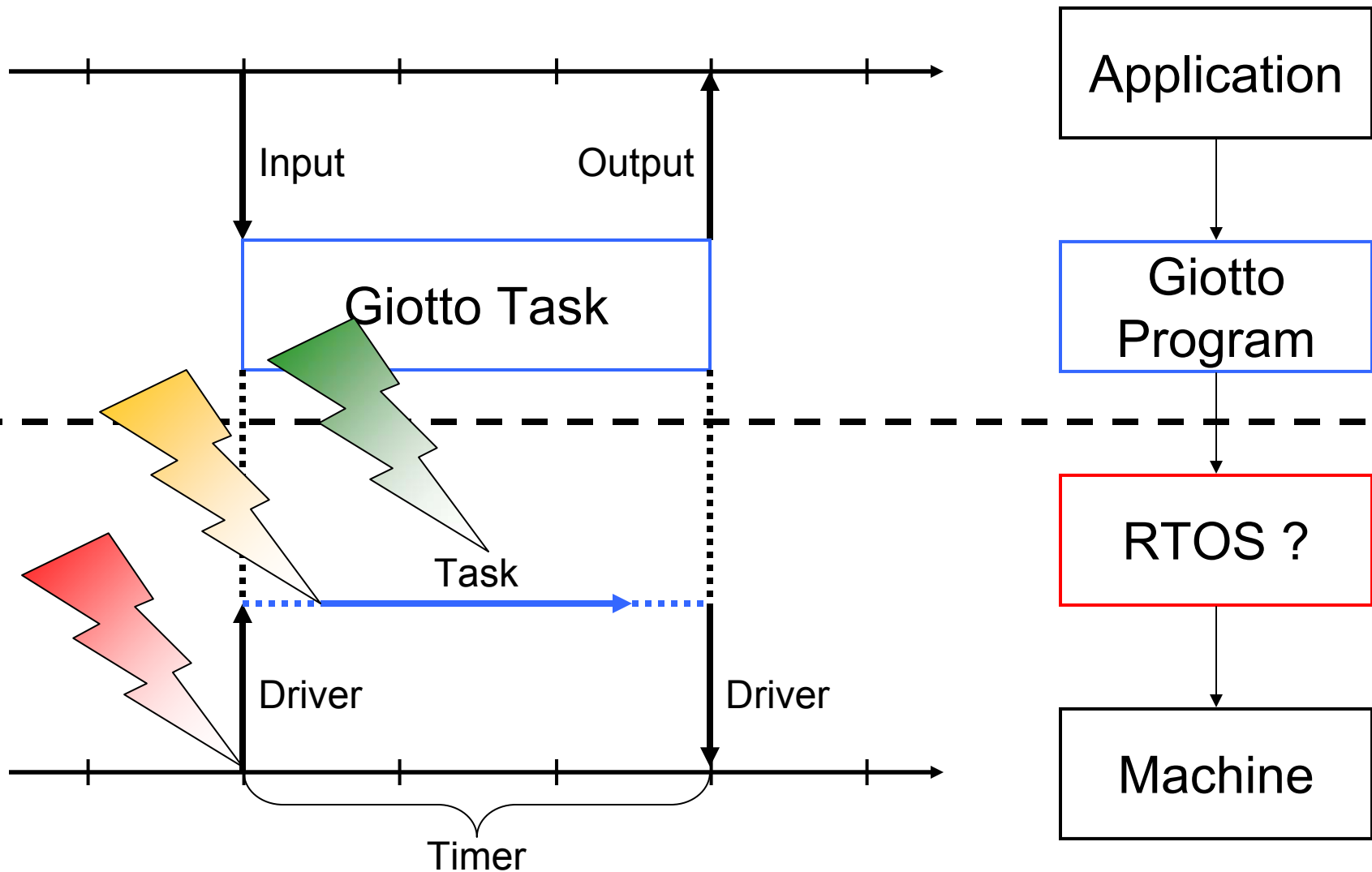


OSEKWorks RTOS, 1 processor (Motorola MPC555 40Mhz)

Operating Systems and Compiler Design: The Embedded Machine and E Code

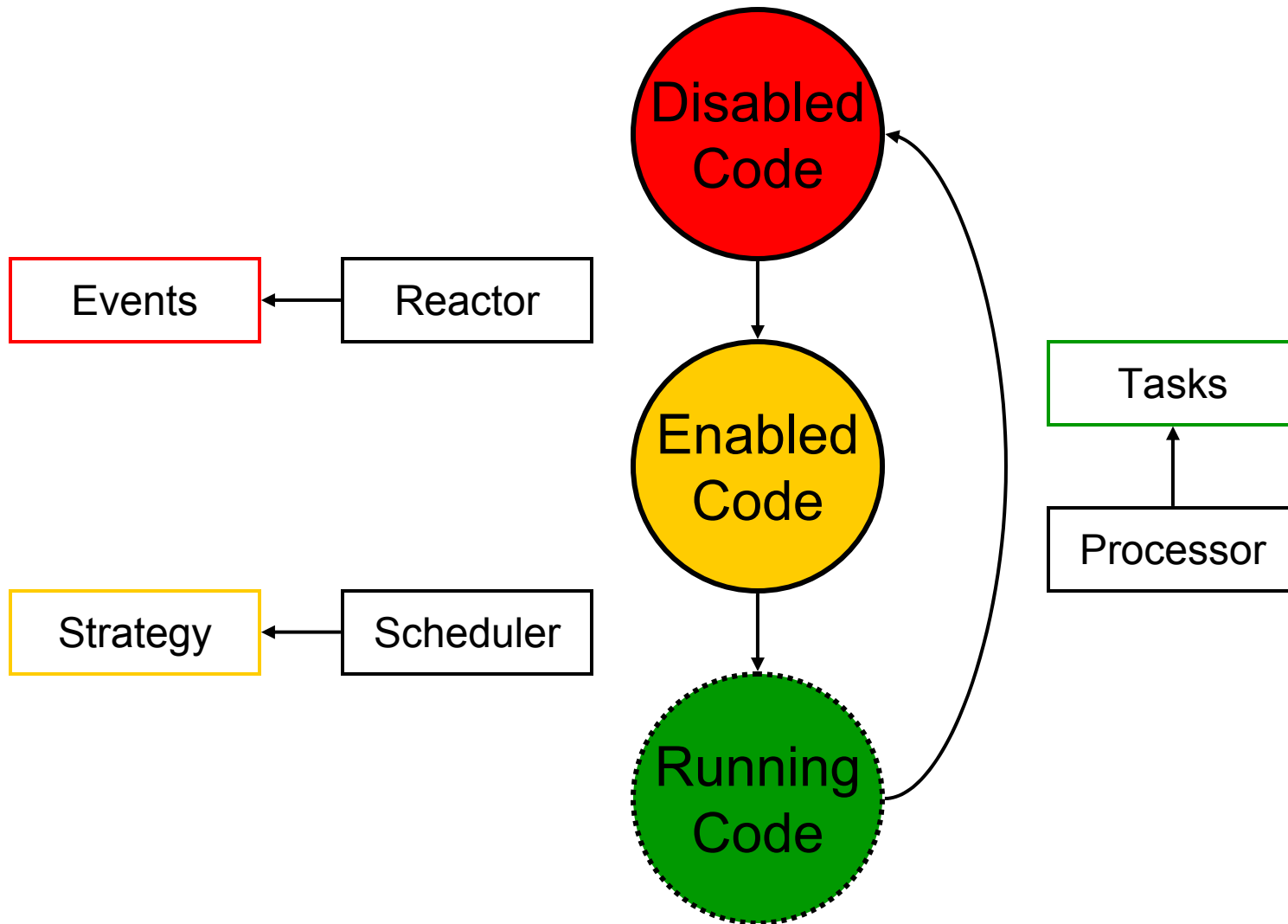


Triggering & Scheduling & Computing

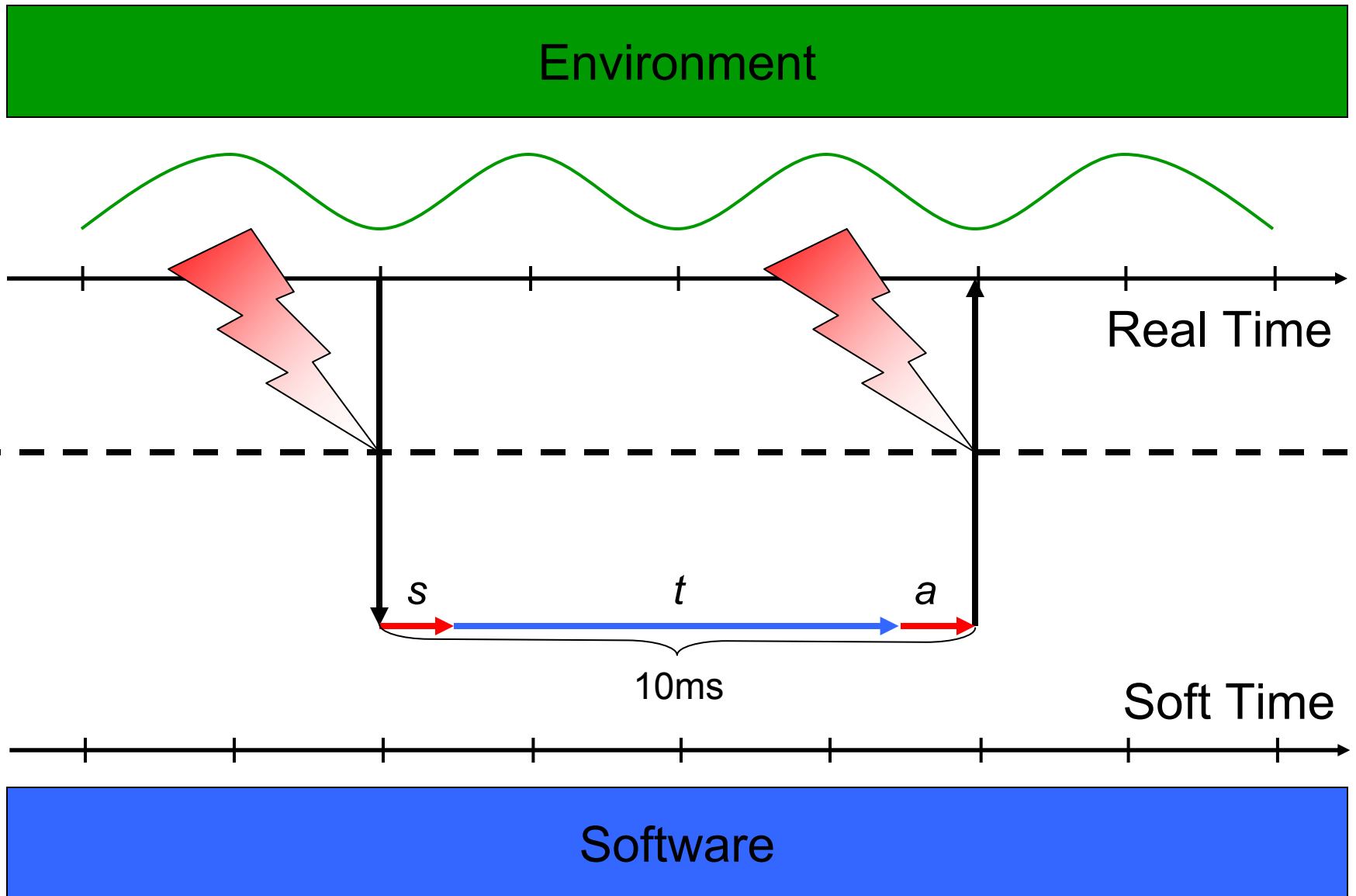


Reactor vs. Scheduler vs. Processor

(Kirsch in the Proceedings of EMSOFT 2002)



Sensing, Computing, Actuating



E Code

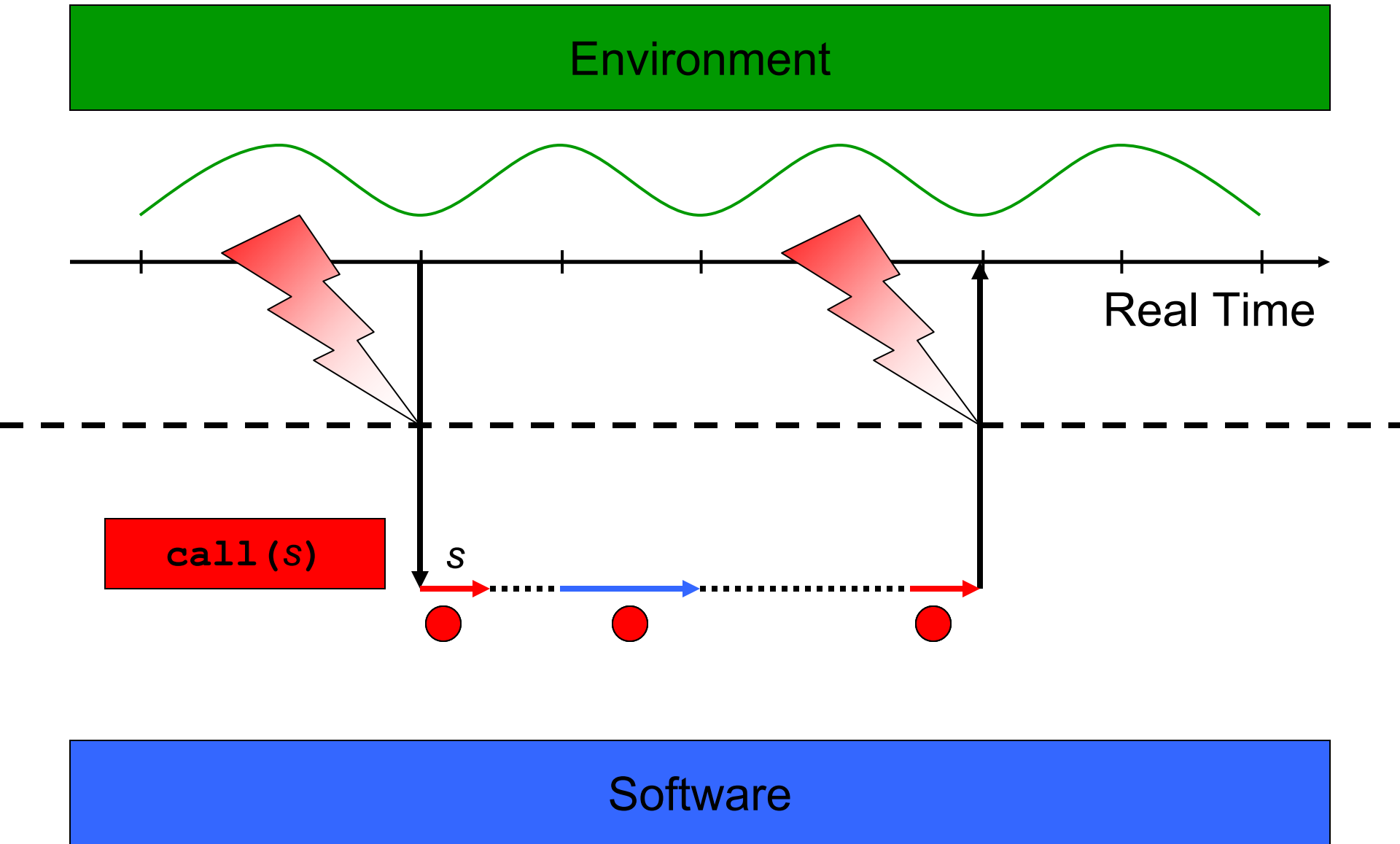
Environment

Real Time

`call (s)`

s

Software



E Code

Environment

Real Time

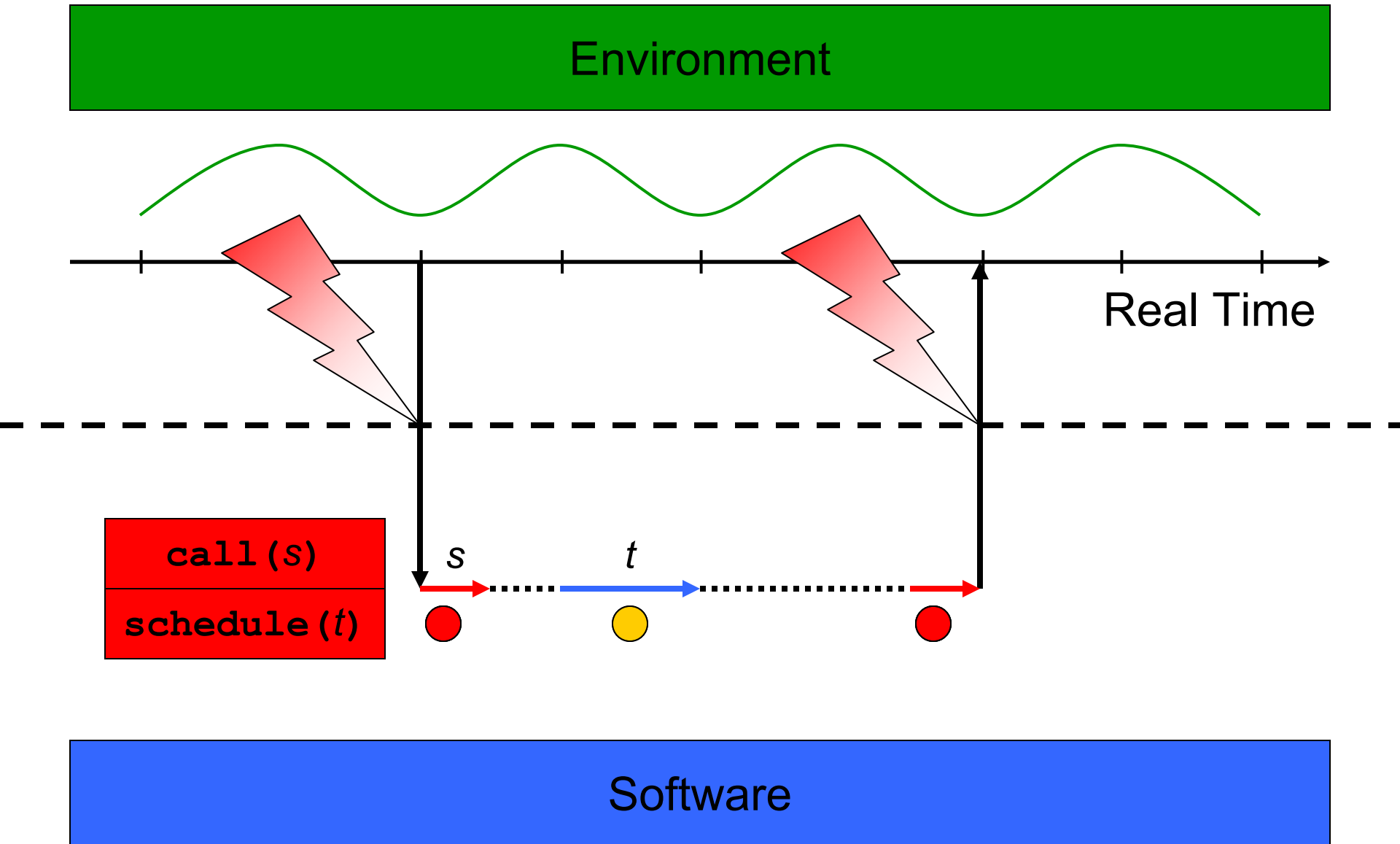
`call (s)`

`schedule (t)`

s

t

Software



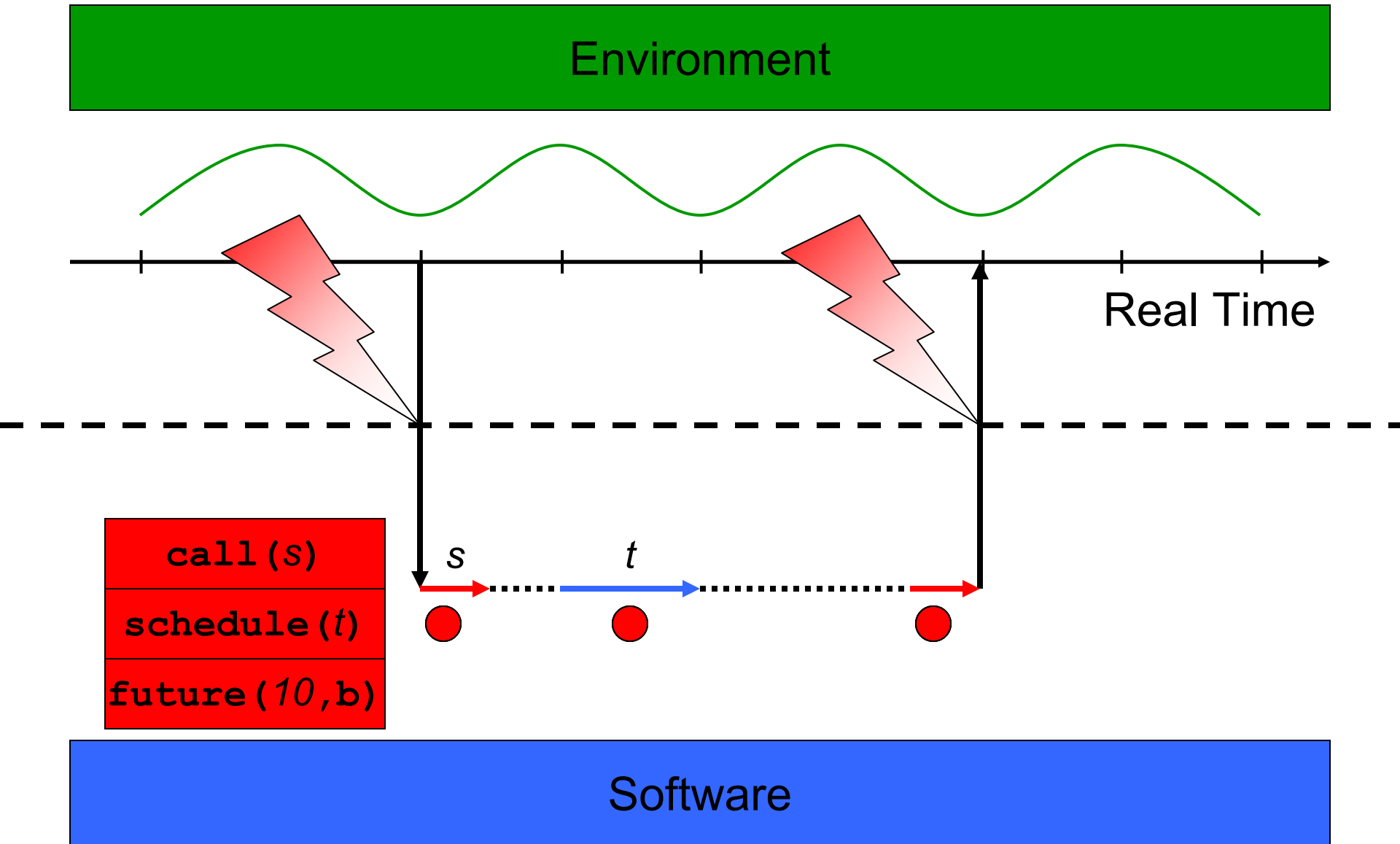
E Code

Environment

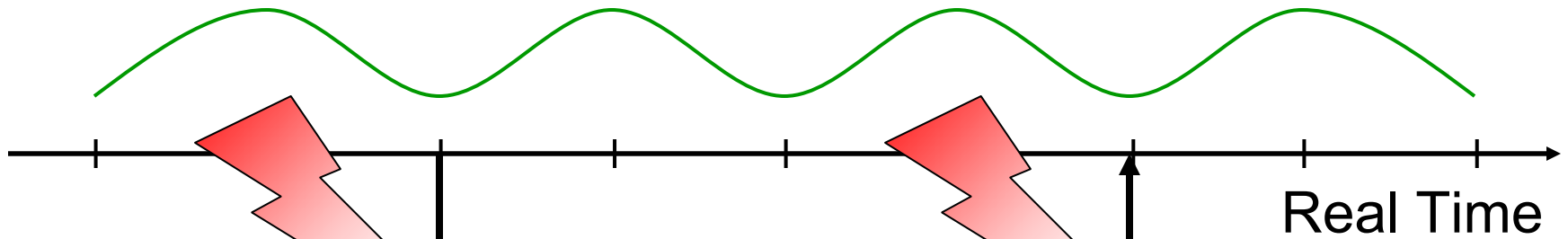
Real Time

<code>call (s)</code>
<code>schedule (t)</code>
<code>future (10, b)</code>

Software

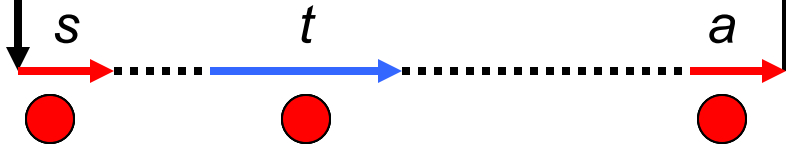


E Code

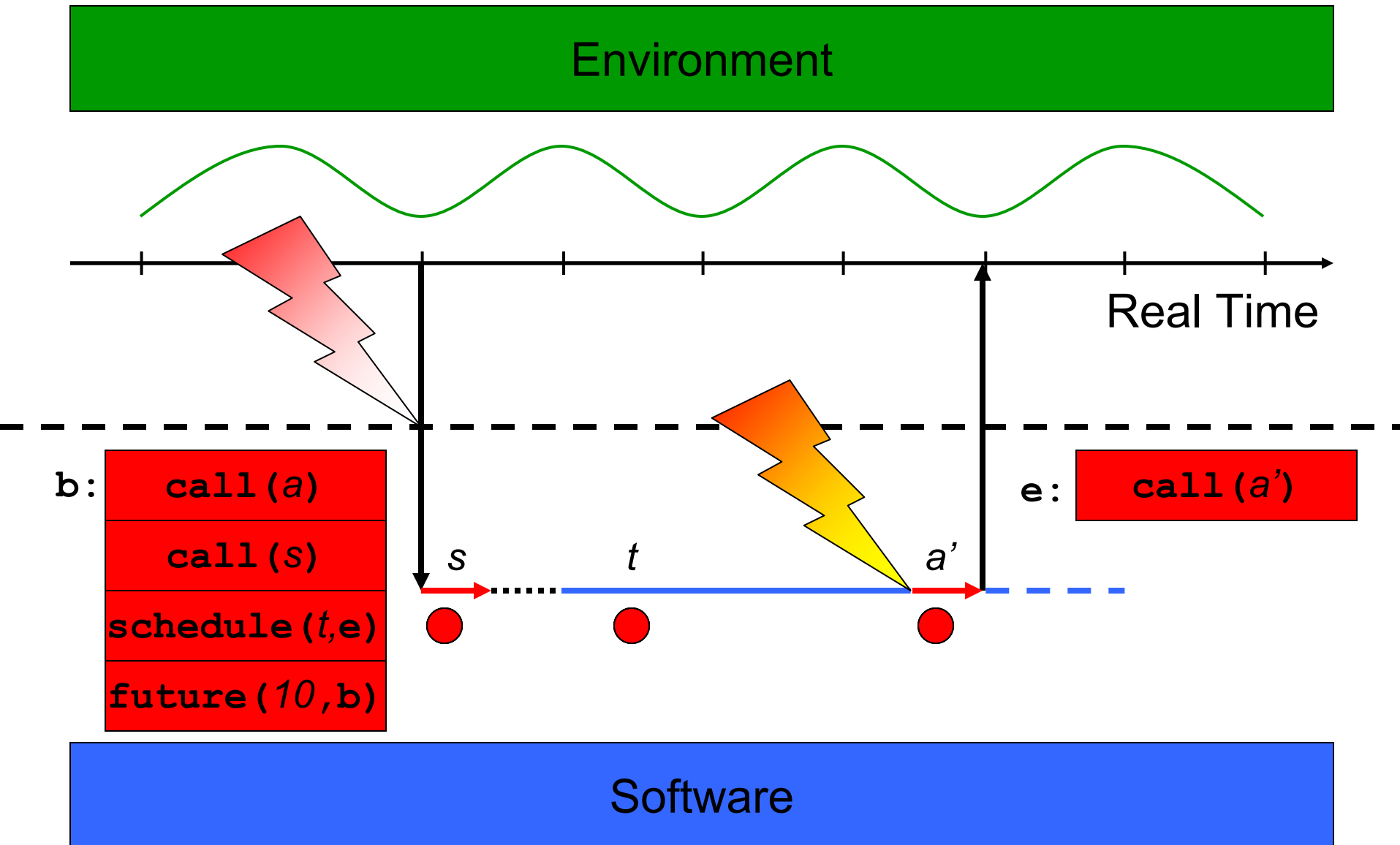


b:

```
call(a)
call(s)
schedule(t)
future(10,b)
```



Time Safety or Runtime Exception



The Embedded Machine

(Henzinger, Kirsch in the Proceedings of PLDI 2002)

The Embedded Machine is a **virtual machine** that **triggers** the execution of software tasks wrt. events

Time-safe, environment-triggered E code is:

- *portable* – **mobile** real-time code
- *predictable* – **deterministic** real-time code
- *composable* – supports:
 - **modular/incremental** compilation
 - **dynamic** linking/patching

The Embedded Machine is available for:

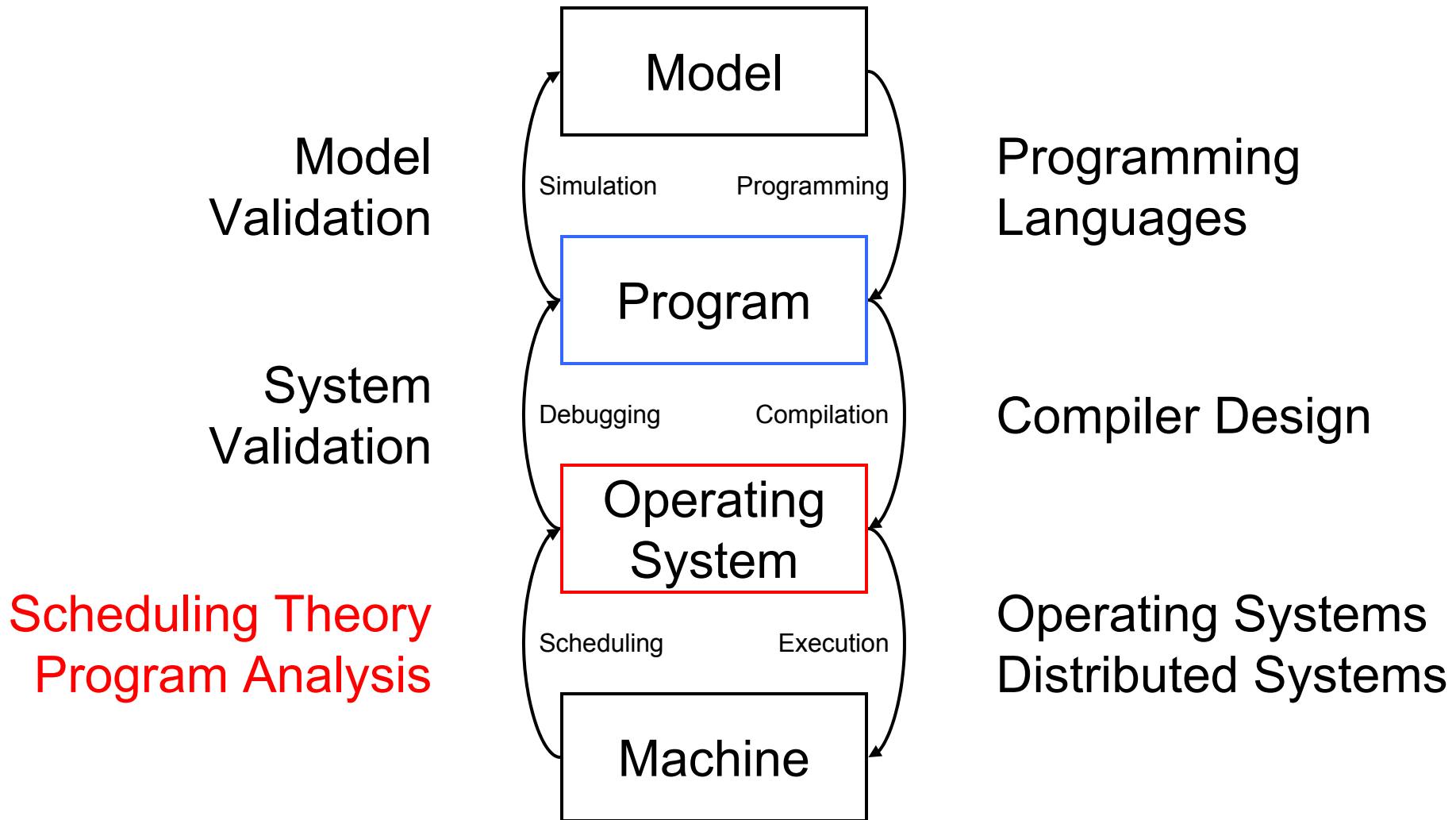
- **Linux**, **Windows**, **OSEKWorks**, **HelyOS**
- **Java** (incl. E code debugger)

Distributed Embedded Machines

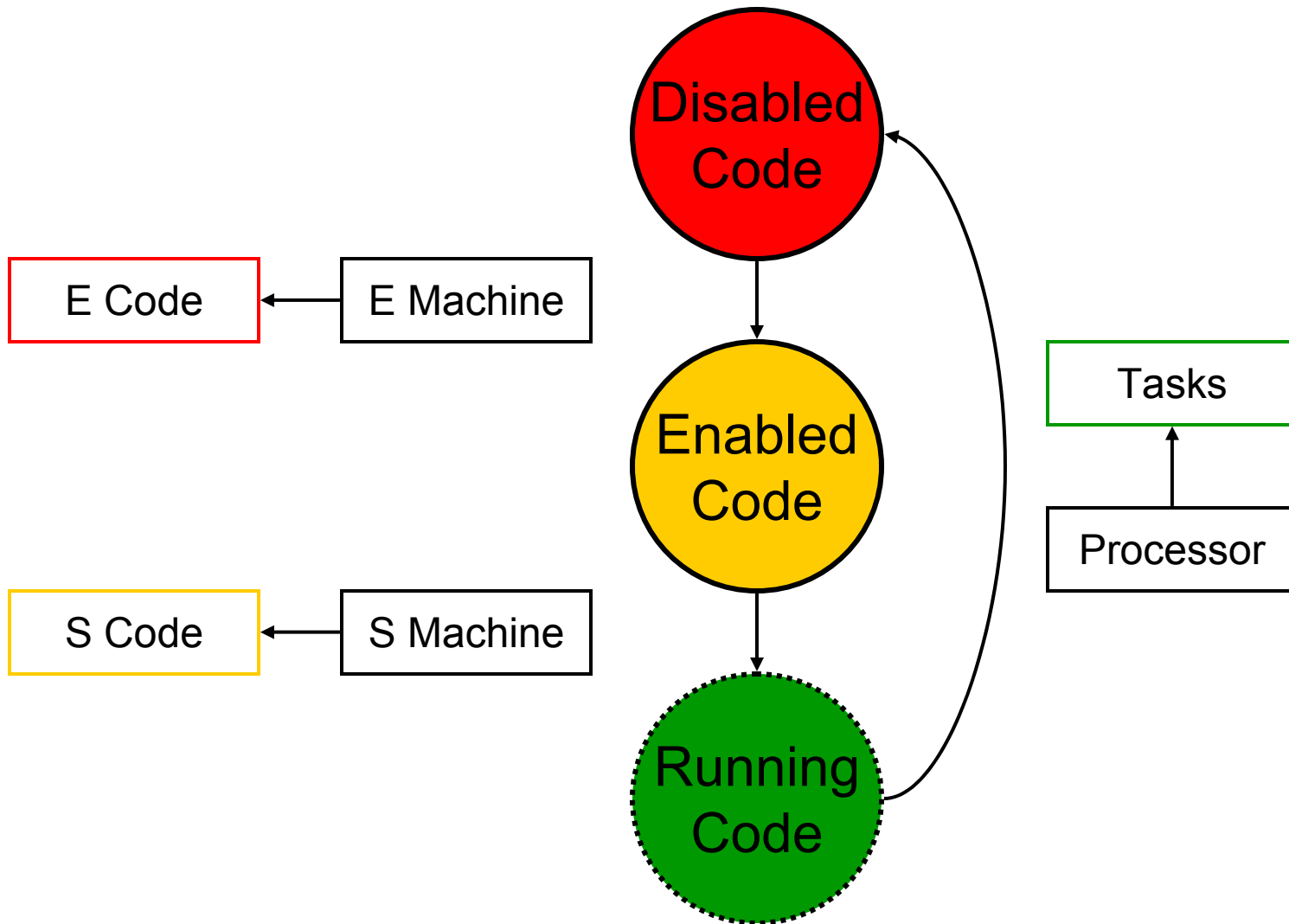
(Part of CHESS, \$13m NSF Project at UC Berkeley)



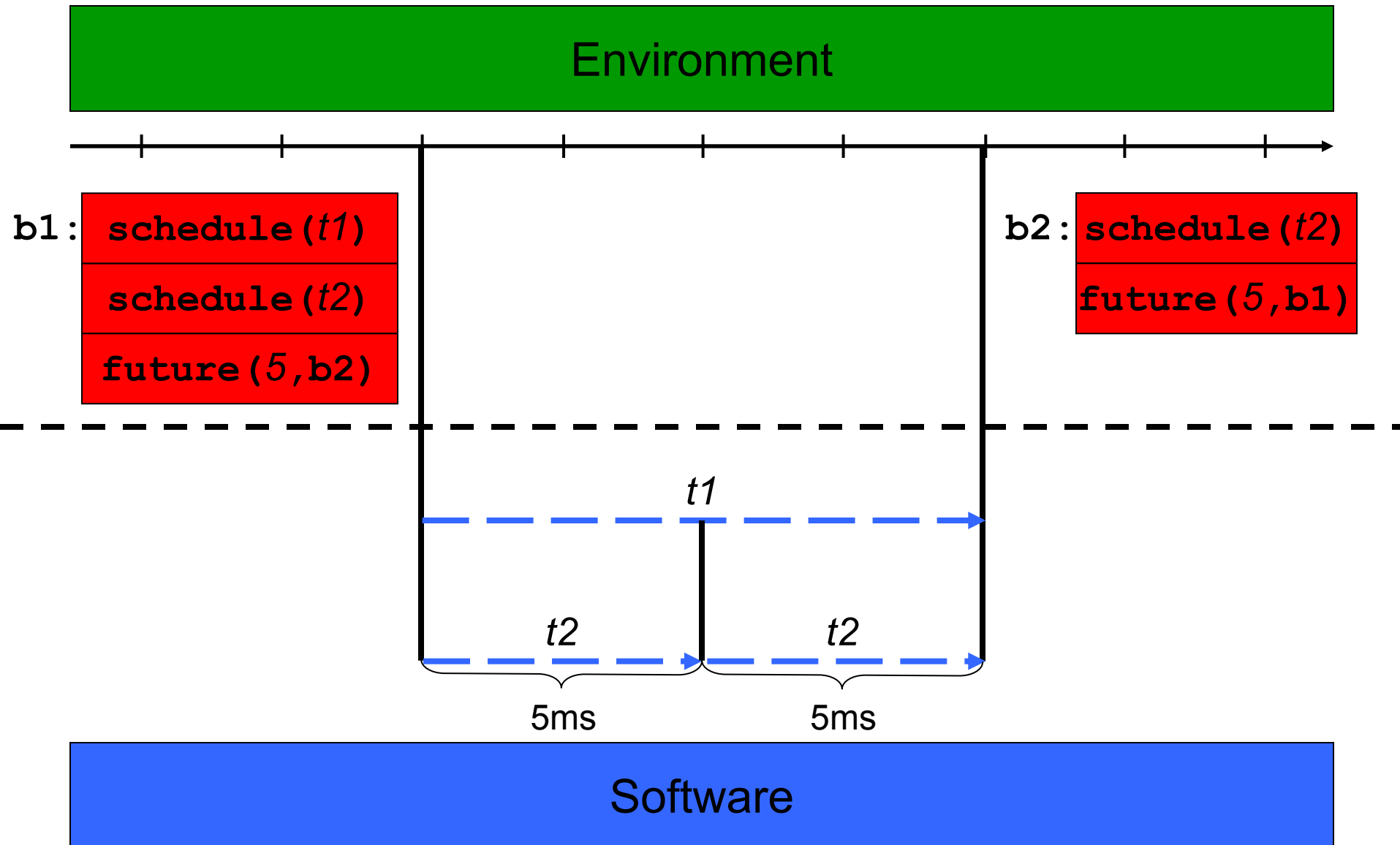
Scheduling Theory and Program Analysis: Schedule-Carrying Code (SCC)



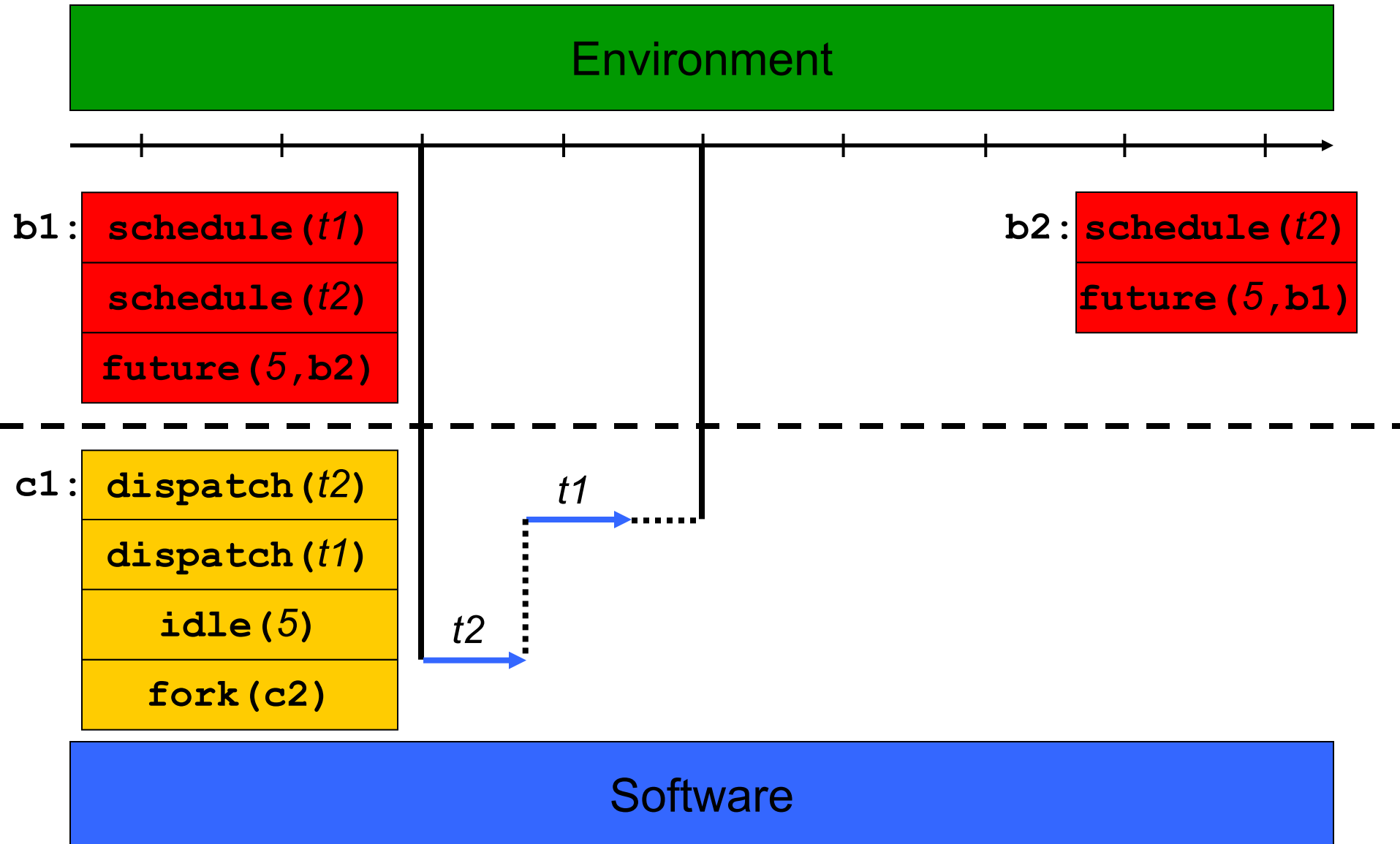
Triggering & Scheduling & Computing



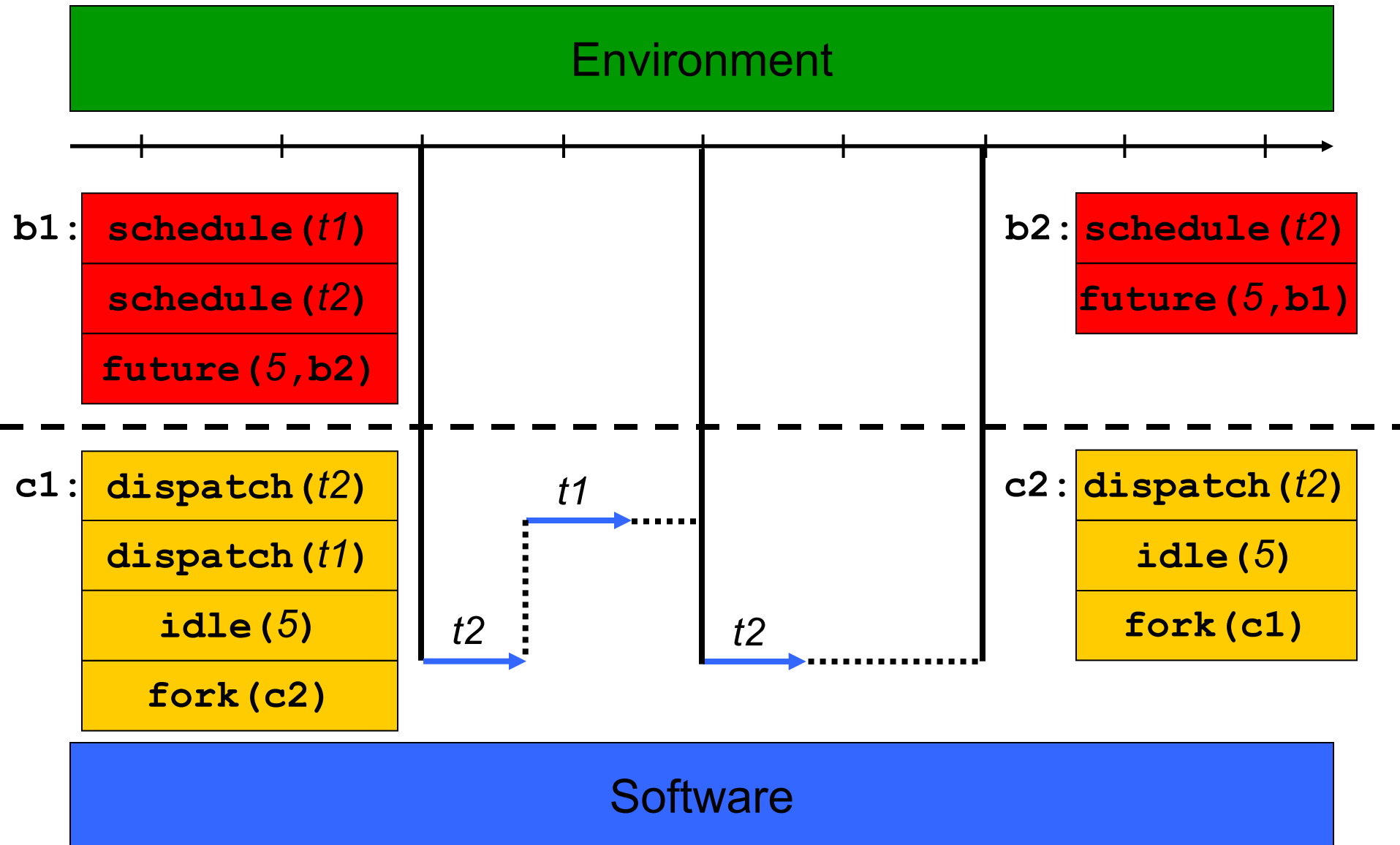
Two Tasks, Different Frequency



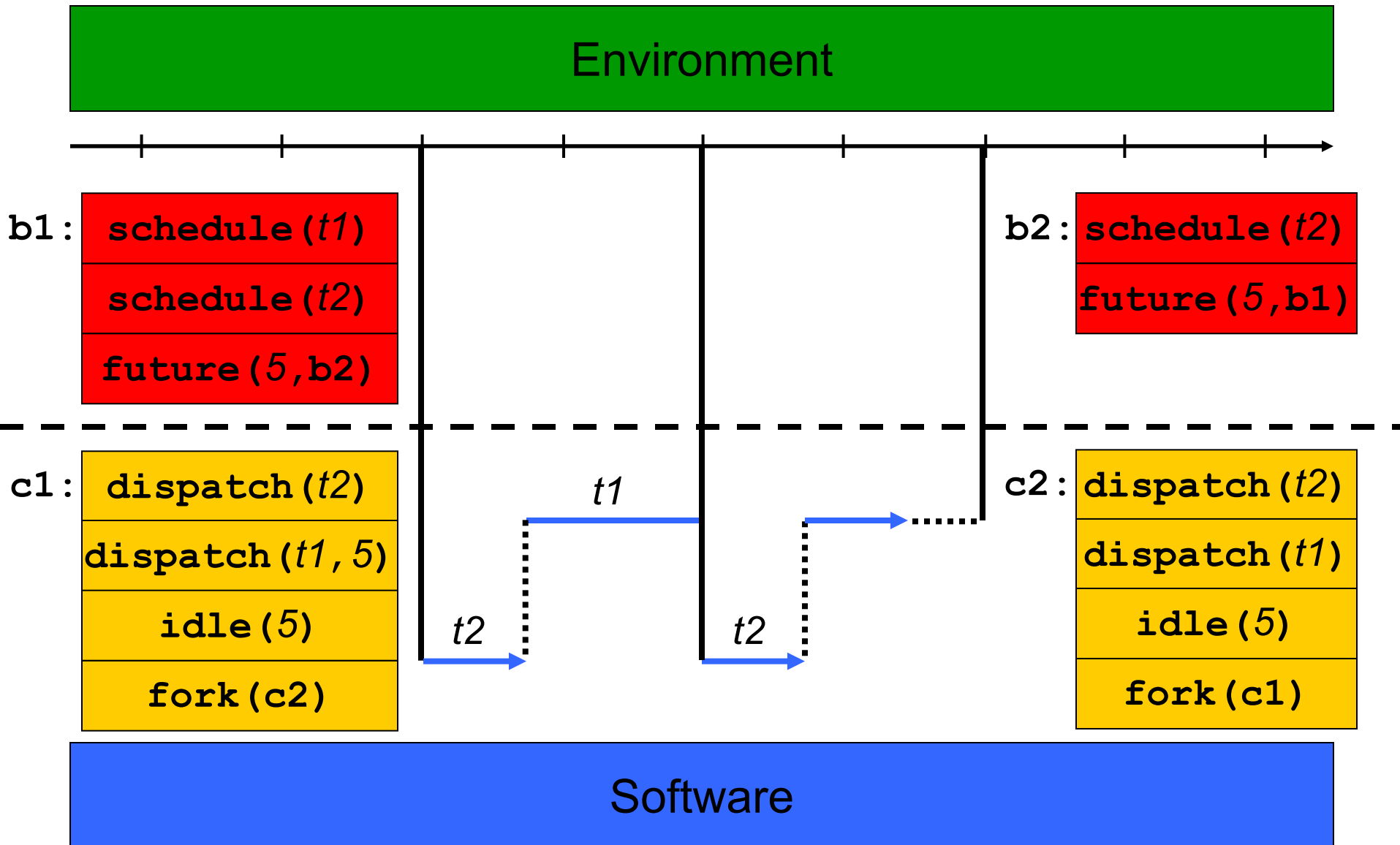
S Code



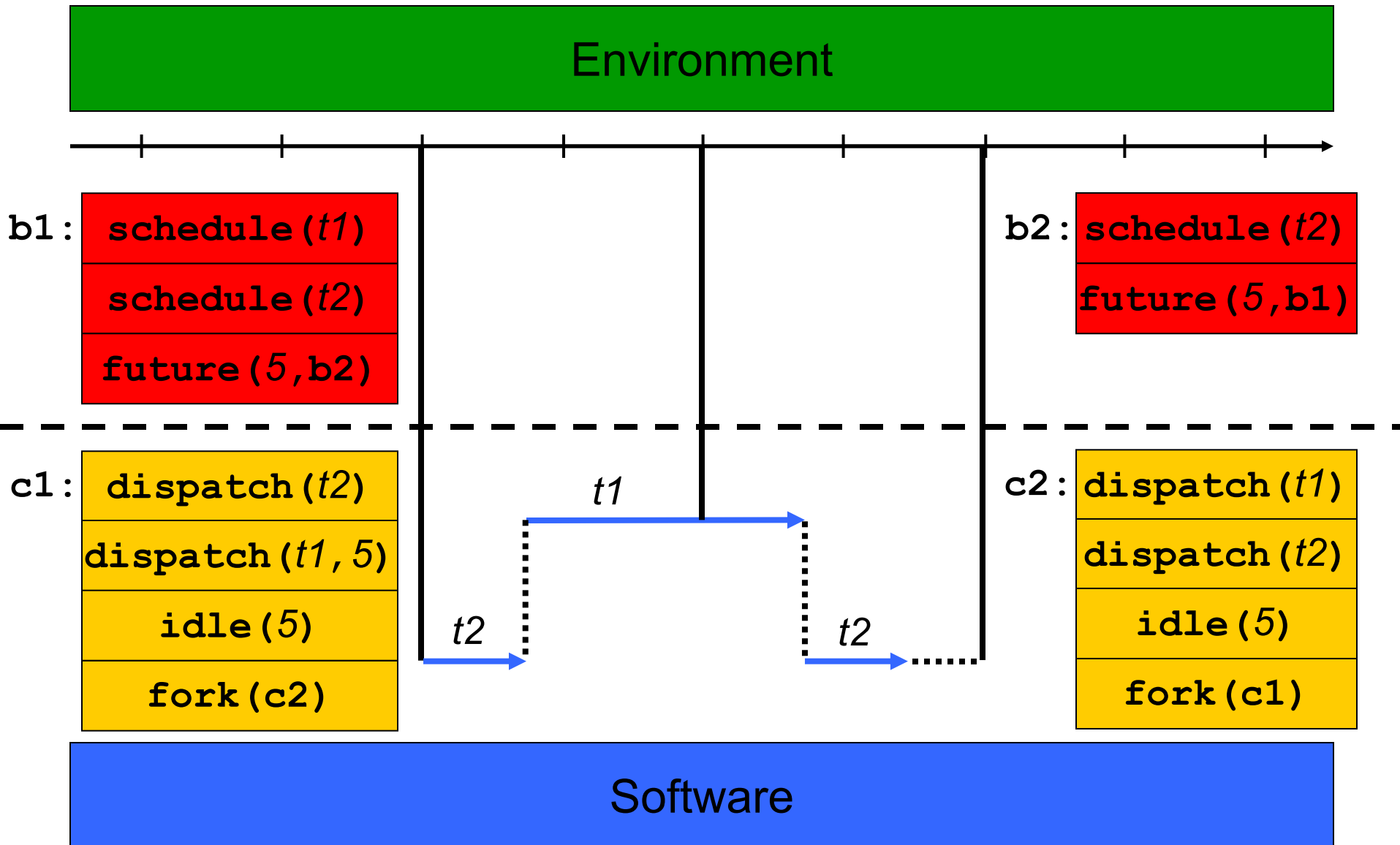
S Code



Preemption



Non-Preemptive Scheduling



Time Safety

Environment

schedule ($t1$)

schedule ($t2$)

future ($5, b2$)

Checking time safety of E code
(set of periodic tasks) with a non-
preemptive scheduler is *NP-hard*

schedule ($t2$)

future ($5, b1$)

Software

Schedule-Carrying Code

(Henzinger, Kirsch, Matic, in Proc. of EMSOFT 2003)

Environment

`schedule (t1)`

`schedule (t2)`

`future (5, b2)`

Checking time safety of E code
(set of periodic tasks) with a non-
preemptive scheduler is *NP-hard*

`schedule (t2)`

`future (5, b1)`

`dispatch (t2)`

`dispatch (t1, 5)`

`idle (5)`

`fork (c2)`

Checking time safety of E code
(set of periodic tasks) + non-
preemptive S code is *linear*
(in E code size)

`dispatch (t1)`

`dispatch (t2)`

`idle (5)`

`fork (c1)`

Software

Time Safety Checking for Embedded Programs:

(Henzinger, Kirsch, Majumdar, Matic in Proc. of EMSOFT 2002)

Scheduling Constraints ↑

Multiprocessor	NP-Complete [Hardness: GJ79]	NP-Hard	NP-Hard	EXPTIME-Complete
Non-Preemptive	NP-Complete [Hardness: JSM91]	NP-Hard	NP-Hard	EXPTIME-Complete
Preemptive	Linear in E Code	Linear in E Code	?	EXPTIME-Complete
	Periodic Tasks	Giotto	?	Arbitrary

→ Triggering Complexity

The Scheduling Machine

The Scheduling Machine is a **virtual machine** that **orders** the execution of software tasks

S code (+ *control-flow* instructions) is:

- *universal* – **any** scheduling strategy
- *verifiable* – **fast** time safety checking
- *distributed* – can schedule:
 - **computation**
 - **communication**

The Scheduling Machine is available for:

- **Java** (incl. S code debugger)

The E and S Machine in Class

Embedded Software Lab 0.2

Go Tick

CI Nn

Complete Complete

0 2000 4000 6000 Real Time

Soft Time

Step	0: schedule(Control)	0: dispatch(Navigation)	Step
Over	1: schedule(Navigation)	1: dispatch(Control)	Over
	2: future(Timer,4,2000)	2: idle(2000)	
	3: return()	3: dispatch(Navigation)	
	4: schedule(Navigation)	4: idle(4000)	
	5: future(Timer,0,2000)	5: fork(0)	
	6: return()	6: return()	

2 Tasks

No Logging

Send E Code

Disconnect

S: pc: -1

Navigation started.
Navigation completed.

S: 6000ms: 4: idle(4000)

S: pc: -1

Summary

